Schur-Weyl Duality for the Clifford Group: Property testing, de Finetti representations, and a robust Hudson theorem

David Gross (Cologne), Sepehr Nezami (Stanford), <u>Michael Walter</u> (Amsterdam)









Schur-Weyl duality

$$(\mathbb{C}^D)^{\otimes t}$$

$$U^{\otimes t} | x_1, \dots, x_t \rangle = U | x_1 \rangle \otimes \dots \otimes U | x_t \rangle$$

$$R_{\pi} | x_1, \dots, x_t \rangle = | x_{\pi^{-1}(1)}, \dots, x_{\pi^{-1}(t)} \rangle$$

Schur-Weyl duality: These actions generate each other's commutant.

$$\frac{U}{U} = -\frac{U}{U}$$

$$\frac{U}{U} = R_{T} U^{2t}$$

Two symmetries that are ubiquituous in quantum information theory:

- i.i.d. quantum information: $[\rho^{\otimes t}, R_{\pi}] = 0$
- eigenvalues, entropies, . . . : $ho \equiv U
 ho U^{\dagger}$
- ightharpoonup randomized constructions: $E_{\mathsf{Haar}}[|\psi\rangle\!\langle\psi|^{\otimes t}]$

See several other talks at this QIP...

Clifford unitaries and stabilizer states

$$\mathbb{C}^D = (\mathbb{C}^d)^{\otimes n}$$

Clifford group: Unitaries U_C such that P Pauli $\Rightarrow U_C P U_C^{\dagger}$ Pauli. For qubits, generated by CNOT, H, P.

Stabilizer states: States of the form $|S\rangle = U_C |0\rangle^{\otimes n}$.

These are important classes of unitaries & states:

- ► QEC, MBQC, topological order, randomized benchmarking, . . .
- ► can be highly entangled, but efficient to represent and compute with
- ▶ 2-design; 3-design for qubits ⇒ efficient random constructions

Motivates a Schur-Weyl duality for the Clifford group!

Our results

"Schur-Weyl duality" for the Clifford group: We characterize precisely which operators commute with $U_C^{\otimes t}$ for all Clifford unitaries U_C .

Fewer unitaries → larger commutant (more than permutations).

Applications:

- Property testing
- ► De Finetti theorems with increased symmetry

$$|S\rangle^{\otimes t}\longleftrightarrow |\psi\rangle^{\otimes t}$$

$$\Psi_s \approx \sum_{S} p_S |S\rangle\langle S|^{\otimes s}$$

► Higher moments of stabilizer states

$$E_S[|S\rangle\!\langle S|^{\otimes t}]$$

- ► t-designs from Clifford orbits
- ► Robust Hudson theorem



Plan:

- Write down permutation action in clever way.
- @ Generalize.
- Prove it!



• Write down permutation action in clever way:

Permutation of t copies of $(\mathbb{C}^d)^{\otimes n}$:

$$R_{\pi} = r_{\pi}^{\otimes n}, \quad r_{\pi} = \sum_{\mathbf{x}} \ket{\pi \mathbf{x}} \langle \mathbf{x} |$$

Here, we think of π as $t \times t$ -permutation matrix, and $|\mathbf{x}\rangle = |x_1, \dots, x_t\rangle$ is computational basis of $(\mathbb{C}^d)^{\otimes t}$.

@ Generalize:

$$R_O = r_O^{\otimes n}, \quad r_O = \sum_{\mathbf{x}} |O\mathbf{x}\rangle \langle \mathbf{x}|$$

Allow all orthogonal and stochastic $t \times t$ -matrices O with entries in \mathbb{F}_d .

For qubits, an example is the 6×6 anti-identity:

$$\overline{\mathrm{id}} = \begin{pmatrix} 0 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 0 & 1 \end{pmatrix},$$

$$R_{\overline{\mathrm{id}}} | \mathbf{x}_1, \dots, \mathbf{x}_6 \rangle = | \mathbf{x}_2 + \dots + \mathbf{x}_6, \dots, \mathbf{x}_1 + \dots + \mathbf{x}_5 \rangle$$

The operator $R_{\overline{id}}$ commutes with $U_C^{\otimes b}$ for every *n*-qubit Clifford unitary.

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The operator $R_{\overline{\mathrm{id}}}$ commutes with $U_C^{\otimes 6}$ for every *n*-qubit Clifford unitary.

Generalize further:

$$R_T = r_T^{\otimes n}, \quad r_T = \sum_{(\boldsymbol{y}, \boldsymbol{x}) \in T} |\boldsymbol{y}\rangle \langle \boldsymbol{x}|$$

Allow all subspaces $T \subseteq \mathbb{F}_d^{2t}$ that are self-dual, i.e. $\mathbf{y} \cdot \mathbf{y} = \mathbf{x} \cdot \mathbf{x}^{\dagger}$ and of dimension t, and contain $\mathbf{1} = (1, \dots, 1)$.

[†]For qubits, require modulo 4 ('doubly even' code).

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Result

For $n \ge t-1$, the operators R_T are $\prod_{k=0}^{t-2} (d^k+1)$ many linearly independent operators that span the commutant of $\{U_C^{\otimes t}\}$.

Sketch of proof: Phase space formalism. Compute cardinalities and compare.



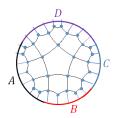
Independent of n! Rich algebraic structure (see paper).

Application 1: Higher moments of stabilizer states

Result (t-th moment)

$$E[|S\rangle\langle S|^{\otimes t}] \simeq \sum_T R_T$$

- ▶ When stabilizer states form *t*-design, reduces to $\sum_{\pi} R_{\pi}$ (Haar average)
- ► Summarizes all previous results on statistical properties
- ▶ ... but also holds for larger t!



We can also write t-th moment as weighted sum of certain CSS codes.

Application 2: Stabilizer testing

Given t copies of an unknown state in $(\mathbb{C}^d)^{\otimes n}$, decide if it is a stabilizer state or ε -far from it.

Idea: Use the anti-identity. Measure POVM element $\frac{1+R_{\overline{id}}}{2}$ on t=6 copies.

Result

Let ψ be a pure state of n qubits. If ψ is a stabilizer state then this accepts always. But if $\max_{\mathcal{S}} |\langle \psi | \mathcal{S} \rangle|^2 \leq 1 - \varepsilon^2$, acceptance probability $\leq 1 - \varepsilon^2/4$.

- ▶ Power of test independent of *n*. Answers q. by Montanaro & de Wolf.
- ► Similar result for qudits & for testing if blackbox unitary is Clifford.

Why does it work? How to implement?

Stabilizer testing using Bell difference sampling

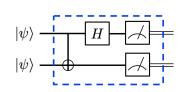
Any state ψ can be expanded in Pauli basis[†]:

$$\psi = \sum_{\mathbf{v}} c_{\psi} P_{\mathbf{v}}$$

- ▶ If pure, then $p_{\psi}(\mathbf{v}) = 2^{n} |c_{\psi}(\mathbf{v})|^{2}$ is a probability distribution.
- ▶ If stabilizer state, then support of p_{ψ} is stabilizer group (up to sign).

Key idea: Sample & verify!

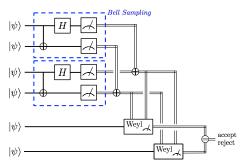
How to sample? If ψ is real, can simply measure in Bell basis $(P_{\mathbf{v}} \otimes I) | \Phi^{+} \rangle$ (Bell sampling; Montanaro, Zhao *et al*).



 $^{{}^{\}dagger}P_{\mathbf{v}} = P_{\nu_1} \otimes \ldots \otimes P_{\nu_n}$ where $P_{00} = I$, $P_{01} = X$, $P_{10} = Z$, $P_{11} = Y$

Stabilizer testing using Bell difference sampling

In general, need to take difference of two Bell measurement outcomes:



- ► Constant-depth circuit, only need coherent two-qubit operations.
- ► Circuit is equivalent to measuring the anti-identity!

Proof of converse uses uncertainty relation.



Application 3: Stabilizer de Finetti theorems

Any tensor power $|\psi\rangle^{\otimes t}$ has S_t -symmetry. De Finetti theorems provide 'partial' converse: If $|\Psi\rangle$ has S_t -symmetry, $\Psi_s \approx \int d\mu(\psi)\psi^{\otimes s}$ for $s \ll t$.

Stabilizer tensor powers have increased symmetry:

$$R_O \ket{S}^{\otimes t} = \ket{S}^{\otimes t}$$
 for all orthogonal and stochastic O

Result

Assume that $|\Psi\rangle \in ((\mathbb{C}^d)^{\otimes n})^{\otimes t}$ has this symmetry, d>2. Then:

$$\|\Psi_s - \sum_{S} p_S |S\rangle\langle S|^{\otimes s}\|_1 \lesssim d^{2n(n+2)} d^{-(t-s)/2}$$

- ▶ Approximation is exponentially good, yet *bona fide* stabilizer states.
- ► Similar to Gaussian de Finetti (Leverrier *et al*). Applications to QKD?

Can reduce symmetry requirements at expense of approximation.

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Application 4: t-designs from Clifford orbits

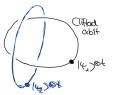
When t>2 or 3 (qubits), stabilizer states fail to be t-design. Yet, hints in the literature that this failure is relatively graceful (Zhu $et\ al$, Nezami-W). E.g., Clifford orbit of non-stabilizer qutrit states can be 3-design!

We prove in general:

Result

For every t, there exists ensemble of N = N(d, t) many fiducial states in $(\mathbb{C}^d)^{\otimes n}$ such that corresponding Clifford orbits form t-design.

- ▶ Number of fiducials does not depend on *n*!
- ► Efficient construction?



Application 5: Robust Hudson theorem

For odd d, every quantum state has a discrete Wigner function:

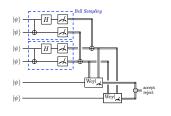
$$W_{
ho}(\mathbf{v}) = d^{-2n} \sum_{\mathbf{w}} e^{-2\pi i [\mathbf{v}, \mathbf{w}]/d} \operatorname{tr}[\rho P_{\mathbf{v}}]$$

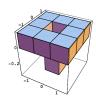


- ▶ Quasi-probability distribution on phase space \mathbb{F}_d^{2n}
- ▶ Discrete Hudson theorem: For pure states, $W_{\psi} \ge 0$ iff ψ stabilizer
- Wigner negativity $\operatorname{sn}(\psi) = \sum_{\mathbf{v}:W_{\rho}(\mathbf{v})<0} |W_{\rho}(\mathbf{v})|$: monotone in resource theory of stabilizer computation; witness for contextuality

Result (Robust Hudson)

There exists a stabilizer state $|S\rangle$ such that $|\langle S|\psi\rangle|^2 \geq 1-9d^2\operatorname{sn}(\psi)$.





- clean algebraic description in terms of self-dual codes
- ► resolve open question in quantum property testing
- ▶ new tools for stabilizer states: moments, de Finetti, Hudson, . . .

Thank you for your attention!