Quantum Brascamp-Lieb Inequalities

Michael Walter



ICMP, August 2021

based on joint work with Mario Berta and David Sutter arXiv:1909.02383

Overview

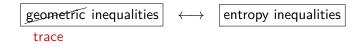
geometric inequalities \longleftrightarrow entropy inequalities

Brascamp-Lieb inequalities have wide range of applications and satisfy beautiful duality. We study a quantum formulation, motivated by the desire to identify new tools to proving entropy inequalities.

Plan for today:

- Introduction
- Quantum BL duality, applications and connections
- Geometric quantum BL inequalities

Overview



Brascamp-Lieb inequalities have wide range of applications and satisfy beautiful duality. We study a quantum formulation, motivated by the desire to identify new tools to proving entropy inequalities.

Plan for today:

- Introduction
- **Q** Quantum BL duality, applications and connections
- Geometric quantum BL inequalities

Classical Brascamp-Lieb inequalities

For $B_k \colon \mathbb{R}^m \twoheadrightarrow \mathbb{R}^{m_k}$ linear, $q_k > 0$, C > 0, an inequality of the form

$$\int_{\mathbb{R}^m} \prod_{k=1}^n |f_k(B_k x)| \, dx \leqslant C \prod_{k=1}^n \|f_k\|_{1/q_k} \qquad \forall f_k$$

This generalizes many classical integral inequalities (Hölder, Young, ...) Many proofs, applications, variations...

• Optimal C can be computed by optimizing over Gaussian f_k . [Lieb]

▶ When is C finite? Fully classified. [Bennett et al]

▶ How to compute *C* efficiently? Still partly open! [Garg et al]

Geometric case: B_k projections s.th. $\sum_{k=1}^n q_k B_k^* B_k = I_m$.

Classical Brascamp-Lieb inequalities

For $B_k: \mathbb{R}^m \twoheadrightarrow \mathbb{R}^{m_k}$ linear, $q_k > 0$, C > 0, an inequality of the form

$$\int_{\mathbb{R}^m} \prod_{k=1}^n |f_k(B_k x)| \, dx \leqslant C \prod_{k=1}^n ||f_k||_{1/q_k} \qquad \forall f_k$$

This generalizes many classical integral inequalities (Hölder, Young, ...) Many proofs, applications, variations...

- Optimal C can be computed by optimizing over Gaussian f_k . [Lieb]
- ▶ When is C finite? Fully classified. [Bennett et al]
- ► How to compute *C* efficiently? Still partly open! [Garg et al]

Geometric case: B_k projections s.th. $\sum_{k=1}^n q_k B_k^* B_k = I_m$.

Duality and entropy

BL inequality is dual to 'subadditivity' inequality for differential entropy:

$$\sum_{k=1}^{n} q_k S(B_k X) \ge S(X) - \log C \qquad \forall \ \mathsf{RV} \ X \ \text{on} \ \mathbb{R}^m$$

Apart from information theoretic interest, equivalence also enables new proof techniques (heat flow). [Carlen-Cordero-Erausquin]

The duality can be generalized to arbitrary channels and relative entropies. Framework includes hypercontractivity, strong data processing, etc. [Liu et al]

Our results: Quantum version of the general duality and applications. In addition, quantum version of the geometric BL inequalities on $L^2(\mathbb{R}^m)$.

Duality and entropy

BL inequality is dual to 'subadditivity' inequality for differential entropy:

$$\sum_{k=1}^{n} q_k S(B_k X) \ge S(X) - \log C \qquad \forall \ \mathsf{RV} \ X \ \text{on} \ \mathbb{R}^m$$

Apart from information theoretic interest, equivalence also enables new proof techniques (heat flow). [Carlen-Cordero-Erausquin]

The duality can be generalized to arbitrary channels and relative entropies. Framework includes hypercontractivity, strong data processing, etc. [Liu et al]

Our results: Quantum version of the general duality and applications. In addition, quantum version of the geometric BL inequalities on $L^2(\mathbb{R}^m)$.

Result: Quantum Brascamp-Lieb Duality

Let $\mathcal{E}_k: L(\mathcal{H}) \to L(\mathcal{H}_k)$ positive & TP, $q_k > 0$, σ , $\sigma_k \succ 0$, C > 0. Then the following are **equivalent**:

$$\sum_{k=1}^{n} q_k D(\mathcal{E}_k(\rho) \| \sigma_k) \leqslant D(\rho \| \sigma) + \log C \qquad \forall \text{ states } \rho$$

and

$$\operatorname{tr} e^{\log \sigma + \sum_{k=1}^{k} \mathcal{E}_{k}^{*}(\log \omega_{k})} \leqslant C \prod_{k=1}^{n} \| e^{\log \omega_{k} + q_{k} \log \sigma_{k}} \|_{1/q_{k}} \quad \forall \ \omega_{k} \succ 0$$

- ► Proof via Legendre: $D(\rho \| \sigma) = \sup_{\omega \succ 0} \{ \operatorname{tr} \rho \log \omega \log \operatorname{tr} e^{\log \omega + \log \sigma} \}$ [Petz]
- Not clear which side looks more intimidating...
- Useful choices: $\sigma_k = \mathcal{E}_k(\sigma)$ or $\sigma_k = I$, $\sigma = I$

Result: Quantum Brascamp-Lieb Duality

Let $\mathcal{E}_k: L(\mathcal{H}) \to L(\mathcal{H}_k)$ positive & TP, $q_k > 0$, σ , $\sigma_k \succ 0$, C > 0. Then the following are **equivalent**:

$$\sum_{k=1}^{n} q_k D(\mathcal{E}_k(\rho) \| \sigma_k) \leqslant D(\rho \| \sigma) + \log C \qquad \forall \text{ states } \rho$$

and

$$\operatorname{tr} e^{\log \sigma + \sum_{k=1}^{k} \mathcal{E}_{k}^{*}(\log \omega_{k})} \leqslant C \prod_{k=1}^{n} \| e^{\log \omega_{k} + q_{k} \log \sigma_{k}} \|_{1/q_{k}} \quad \forall \ \omega_{k} \succ 0$$

- ► Proof via Legendre: $D(\rho \| \sigma) = \sup_{\omega \succ 0} \{ \operatorname{tr} \rho \log \omega \log \operatorname{tr} e^{\log \omega + \log \sigma} \}$ [Petz]
- Not clear which side looks more intimidating...
- Useful choices: $\sigma_k = \mathcal{E}_k(\sigma)$ or $\sigma_k = I$, $\sigma = I$

Without side information

[Carlen-Lieb]

When specializing to $\sigma_k = I$, $\sigma = I$, recover equivalence between

$$\sum_{k=1}^{n} q_k S(\mathcal{E}_k(\rho)) \ge S(\rho) - \log C \qquad \forall \text{ states } \rho$$

and

$$\operatorname{tr} e^{\sum_{k=1}^{k} \mathcal{E}_{k}^{*}(\log \omega_{k})} \leqslant C \prod_{k=1}^{n} \|\omega_{k}\|_{1/q_{k}} \qquad \forall \ \omega_{k} \succ 0$$

For example, can prove uncertainty relations via trace inequalities, as pioneered by Frank-Lieb:

- Maassen-Uffink: $S(X) + S(Z) \ge S(\rho) + 1$ via Golden-Thompson
- Six-state [Coles et al]: $S(X) + S(Y) + S(Z) \ge S(\rho) + 2$ via Lieb 3-matrix

Without side information

[Carlen-Lieb]

When specializing to $\sigma_k = I$, $\sigma = I$, recover equivalence between

$$\sum_{k=1}^{n} q_k S(\mathcal{E}_k(\rho)) \ge S(\rho) - \log C \qquad \forall \text{ states } \rho$$

and

$$\operatorname{tr} e^{\sum_{k=1}^{k} \mathcal{E}_{k}^{*}(\log \omega_{k})} \leqslant C \prod_{k=1}^{n} \|\omega_{k}\|_{1/q_{k}} \qquad \forall \ \omega_{k} \succ 0$$

For example, can prove uncertainty relations via trace inequalities, as pioneered by Frank-Lieb:

- Maassen-Uffink: $S(X) + S(Z) \ge S(\rho) + 1$ via Golden-Thompson
- ► Six-state [Coles et al]: $S(X) + S(Y) + S(Z) \ge S(\rho) + 2$ via Lieb 3-matrix

Applications and questions

- Can we prove new uncertainty relations involving multiple measurements (and even general quantum channels)? N-matrix GT?
- Strong data-processing inequalities fall into the framework:

 $D(\mathcal{E}(\rho) \| \mathcal{E}(\sigma)) \leqslant \eta D(\rho \| \sigma) \qquad \forall \rho$

Tensorization holds classically, but fails quantumly:

 $(\mathcal{E}, \mathcal{C}) \& (\mathcal{E}', \mathcal{C}') \not\Rightarrow (\mathcal{E} \otimes \mathcal{E}', \mathcal{C} \cdot \mathcal{C}')$

Examples include non-additivity of minimal output entropy. Useful?

- Computational complexity of testing validity of (families of) BL ineqs?
- Relation to works by Carlen-Maas?

Back to geometry...

Recall the classical Brascamp-Lieb inequalities in the geometric case:

$$\sum_{k=1}^{n} q_k S(P_k X) \ge S(X) \qquad \forall \ \mathsf{RV} \ X \ \mathsf{on} \ \mathbb{R}^m$$

with P_k projections onto subspaces $V_k \subseteq \mathbb{R}^m$ s.th. $\sum_{k=1}^n q_k P_k = I_m$.

How can we formulate a quantum version? For any subspace $V \subseteq \mathbb{R}^m$,

$$L^{2}(\mathbb{R}^{m}) = L^{2}(V \oplus V^{\perp}) = L^{2}(V) \otimes L^{2}(V^{\perp})$$

hence can define reduced state ρ_V for any state ρ on $L^2(\mathbb{R}^m)$.

This generalizes the usual partial trace. In general, can interpret as state of subset of modes after subjecting ρ to network of beamsplitters.

Back to geometry...

Recall the classical Brascamp-Lieb inequalities in the geometric case:

$$\sum_{k=1}^{n} q_k S(P_k X) \ge S(X) \qquad \forall \ \mathsf{RV} \ X \ \mathsf{on} \ \mathbb{R}^m$$

with P_k projections onto subspaces $V_k \subseteq \mathbb{R}^m$ s.th. $\sum_{k=1}^n q_k P_k = I_m$.

How can we formulate a quantum version? For any subspace $V \subseteq \mathbb{R}^m$,

$$L^{2}(\mathbb{R}^{m}) = L^{2}(V \oplus V^{\perp}) = L^{2}(V) \otimes L^{2}(V^{\perp})$$

hence can define reduced state ρ_V for any state ρ on $L^2(\mathbb{R}^m)$.

This generalizes the usual partial trace. In general, can interpret as state of subset of modes after subjecting ρ to network of beamsplitters.

Result: Geometric Quantum Brascamp-Lieb Inequality

Theorem

Let P_k projections onto subspaces $V_k \subseteq \mathbb{R}^m$ s.th. $\sum_{k=1}^n q_k P_k = I_m$. Then, for all states ρ on $L^2(\mathbb{R}^m)$ with finite first and second moments: $\sum_{k=1}^n q_k S(\rho_{V_k}) \ge S(\rho)$

- ► For coordinate subspaces recover quantum Shearer inequality. [Carlen-Lieb]
- But already nontrivial for "Mercedes star" configuration in \mathbb{R}^2 :

- ► Also holds conditioned on side information. [Ligthard]
- ▶ Can generate more ineqs. via Gaussian unitaries: $Sp_{2m} \frown L^2(\mathbb{R}^m)...$

Sketch of proof

$$\sum_{k=1}^n q_k \, S(\rho_{V_k}) \geqslant S(\rho)$$

Implement classical proof strategy of Carlen–Cordero-Erausquin using quantum heat flow of König-Smith: cf. [De Palma–Trevisan]

$$rac{d}{dt}
ho = -\sum_{j=1}^m [Q_j, [Q_j,
ho]] + [P_j, [P_j,
ho]]$$

Asymptotic scaling of entropy: $S(\rho_V(t)) \sim \dim V \log t$

• Inequality holds at $t = \infty$ if $\sum_k q_k \dim V_k \ge m$.

Quantum de Bruijn identity: $\frac{d}{dt}S(\rho) = J(\rho)$, a Fisher information.

• Can prove reverse inequality for Fisher information if $\sum_k q_k P_k \leq I_m$:

$$\sum_{k=1}^n q_k J(\rho_{V_k}) \leqslant J(\rho)$$

Sketch of proof

$$\sum_{k=1}^n q_k \, S(\rho_{V_k}) \geqslant S(\rho)$$

Implement classical proof strategy of Carlen–Cordero-Erausquin using quantum heat flow of König-Smith: cf. [De Palma–Trevisan]

$$\frac{d}{dt}\rho = -\sum_{j=1}^{m} [Q_j, [Q_j, \rho]] + [P_j, [P_j, \rho]]$$

Asymptotic scaling of entropy: $S(\rho_V(t)) \sim \dim V \log t$

• Inequality holds at $t = \infty$ if $\sum_k q_k \dim V_k \ge m$.

Quantum de Bruijn identity: $\frac{d}{dt}S(\rho) = J(\rho)$, a Fisher information.

• Can prove reverse inequality for Fisher information if $\sum_k q_k P_k \leq I_m$:

$$\sum_{k=1}^n q_k J(\rho_{V_k}) \leqslant J(\rho)$$

Gaussian BL beyond the geometric case

There is a natural action of Sp_{2m} on $L^2(\mathbb{R}^m)$ by Gaussian unitaries. Any symplectic matrix $B \in \mathbb{R}^{2m' \times 2m}$ determines subsystem of m' modes, so we can define reduced state ρ_B on $L^2(\mathbb{R}^{m'})$ for any state ρ on $L^2(\mathbb{R}^m)$.

This notion generalizes the reduced state ρ_V for subspaces $V \subseteq \mathbb{R}^m$ and leads naturally to the following class of Gaussian quantum BL inequalities:

$$\sum_{k=1}^n q_k S(\rho_{B_k}) \geqslant S(\rho) + c$$

where the $B_k \in \mathbb{R}^{2m_k \times 2m}$ symplectic matrices. When does it hold?

Recent result (De Palma–Trevisan): Assuming $\sum_{k=1}^{n} q_k m_k = m$, inequality holds for all quantum states iff holds for all probability densities!

Gaussian BL beyond the geometric case

There is a natural action of Sp_{2m} on $L^2(\mathbb{R}^m)$ by Gaussian unitaries. Any symplectic matrix $B \in \mathbb{R}^{2m' \times 2m}$ determines subsystem of m' modes, so we can define reduced state ρ_B on $L^2(\mathbb{R}^{m'})$ for any state ρ on $L^2(\mathbb{R}^m)$.

This notion generalizes the reduced state ρ_V for subspaces $V \subseteq \mathbb{R}^m$ and leads naturally to the following class of Gaussian quantum BL inequalities:

$$\sum_{k=1}^n q_k S(
ho_{B_k}) \geqslant S(
ho) + c$$

where the $B_k \in \mathbb{R}^{2m_k \times 2m}$ symplectic matrices. When does it hold?

Recent result (De Palma–Trevisan): Assuming $\sum_{k=1}^{n} q_k m_k = m$, inequality holds for all quantum states iff holds for all probability densities!

Gaussian BL beyond the geometric case

There is a natural action of Sp_{2m} on $L^2(\mathbb{R}^m)$ by Gaussian unitaries. Any symplectic matrix $B \in \mathbb{R}^{2m' \times 2m}$ determines subsystem of m' modes, so we can define reduced state ρ_B on $L^2(\mathbb{R}^{m'})$ for any state ρ on $L^2(\mathbb{R}^m)$.

This notion generalizes the reduced state ρ_V for subspaces $V \subseteq \mathbb{R}^m$ and leads naturally to the following class of Gaussian quantum BL inequalities:

$$\sum_{k=1}^n q_k S(\rho_{B_k}) \geqslant S(\rho) + c$$

where the $B_k \in \mathbb{R}^{2m_k \times 2m}$ symplectic matrices. When does it hold?

Recent result (De Palma–Trevisan): Assuming $\sum_{k=1}^{n} q_k m_k = m$, inequality holds for all quantum states iff holds for all probability densities!

- Also holds conditioned on side information.
- ► Can also include "classical" outputs (= quadrature measurements)
- Proof again based on quantum heat flow strategy!

Outlook

. . .

$$\begin{array}{c|c} \hline \text{trace inequalities} & \xleftarrow{BL} \\ \xrightarrow{duality} & entropy inequalities \end{array}$$

Duality between quantum relative entropy inequalities and trace inequalities. Unifying framework to tackle information theoretic questions. New family of geometric quantum Brascamp-Lieb inequalities.

Many exciting directions:

- ▶ Uncertainty relations from *n*-matrix GT?
- Sufficient conditions for tensorization?
- Applications of new trace inequalities?
- Other applications of quantum heat flow?

Thank you for your attention!