

Distance Measures & Intro to Q Channels

Invariant under $\Gamma \rightarrow V\Gamma W$
for some V, W

Operator norms: For arbitrary $M \in L(\mathcal{H}, \mathcal{K})$ w/ singular values $\{s_i\}$

* trace norm: $\|M\|_1 = \sum_i s_i = \text{tr}[\sqrt{M^\dagger M}]$ ← Square root of PSD operator

* Frobenius norm: $\|M\|_2 = \sqrt{\sum_i s_i^2} = \sqrt{\text{tr}[M^\dagger M]}$ ← induced by Hilbert-Schmidt inner product

* operator norm: $\|M\|_\infty = \max_i s_i = \max_{\|\phi\|=1} \|M|\phi\rangle\|$ ($\langle M, N \rangle := \text{tr}[M^\dagger N]$)

USEFUL: submultiplicative

& $|\text{tr}[AB]| \leq \begin{cases} \|A\|_2 \cdot \|B\|_2 & \text{Cauchy-Schwarz ieq} \\ \|A\|_1 \cdot \|B\|_\infty & \text{Hölder ieq} \end{cases}$

$$\|A\|_1 = \max_{\|B\|_\infty \leq 1} |\text{tr}[AB]| = \max_{U \text{ unitary}} |\text{tr}[AU]| \text{ for } A \in L(\mathcal{H})$$

\geq Hölder
 \geq clear
 $\geq \|A\|_1$ by choosing $U|e_i\rangle = |g_i\rangle$

(Normalized) trace distance between $\rho, \sigma \in D(\mathcal{H})$:

$$T(\rho, \sigma) = \frac{1}{2} \|\rho - \sigma\|_1 \in [0, 1]$$

* $T(\rho, \sigma) = \sqrt{1 - |\langle \Phi | \Psi \rangle|^2}$ for $\rho = |\Psi\rangle\langle\Psi|$, $\sigma = |\Phi\rangle\langle\Phi|$ pure HLW 1

* $T(\rho, \sigma) = \max_{0 \leq Q \leq I} \text{tr}[Q(\rho - \sigma)]$ HLW 2 * nice interpretation! (Helmsdon's theorem)

Now comes a strange definition...

Fidelity between $\rho, \sigma \in D(\mathcal{H})$:

$$F(\rho, \sigma) := \|\sqrt{\rho} \sqrt{\sigma}\|_1 = \text{tr} \sqrt{\rho \sigma \rho}$$

NOT Hermit!

* If $\rho = |\Psi\rangle\langle\Psi|$ pure: $F(\rho, \sigma) \stackrel{\text{use}}{=} \sqrt{\langle \Psi | \sigma | \Psi \rangle} \stackrel{\text{use}}{=} |\langle \Psi | \Phi \rangle|$
 $\sqrt{\rho} = \rho$ ↑ if also $\sigma = |\Phi\rangle\langle\Phi|$ pure

* $F(\rho, \sigma) = F(\sigma, \rho)$

* $0 \leq F(\rho, \sigma) \leq 1, F(\rho, \sigma) = \begin{cases} 1 & \text{if } \rho = \sigma \\ 0 & \text{if } \rho\sigma = 0 \end{cases}$

Similarity measure!

Uhlmann's Theorem: If ρ_A, σ_A have purifications on $\mathcal{H}_A \otimes \mathcal{H}_B$:
 $F(\rho_A, \sigma_A) = \max \{ |\langle \Psi_{AB} | \Phi_{AB} \rangle| : |\Psi_{AB}\rangle, |\Phi_{AB}\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B \}$
 purifications of ρ_A, σ_A

NB: RHS = $\max \{ |\langle \Psi_{AB}^{\text{fixed}} | I_A \otimes U_B | \Phi_{AB}^{\text{fixed}} \rangle| : U_B \text{ unitary on } \mathcal{H}_B \}$

Pf: For simplicity, assume $\mathcal{H}_A = \mathcal{H}_B$. Take

$|\Psi_{AB}^{\text{std}}\rangle = (\sqrt{\rho_A} \otimes I_B) \sum_x |x\rangle|x\rangle$ & $|\Phi_{AB}^{\text{std}}\rangle = (\sqrt{\sigma_A} \otimes I_B) \sum_x |x\rangle|x\rangle$

$\Rightarrow |\langle \Psi_{AB}^{\text{std}} | I_A \otimes U_B | \Phi_{AB}^{\text{std}} \rangle| = \sum_{x,y} (\langle x | \otimes \langle x |) (\sqrt{\rho_A} \sqrt{\sigma_A} \otimes U_B) |y\rangle |y\rangle$

$= \sum_{x,y} \langle x | \sqrt{\rho_A} \sqrt{\sigma_A} |y\rangle \underbrace{\langle x | U_B |y\rangle}_{=\langle y | U_B^T |x\rangle} = \text{tr} [\sqrt{\rho_A} \sqrt{\sigma_A} U_B^T]$
 arbitrary unitary

$\Rightarrow \max_{U_B} \|\sqrt{\rho_A} \sqrt{\sigma_A}\|$

□

What if $\mathcal{H}_A \neq \mathcal{H}_B$? See lecture notes!

Properties of the fidelity:

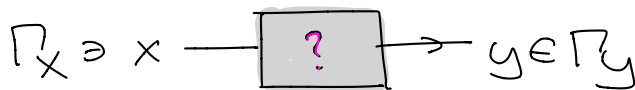
* **Monotonicity:** $F(\rho_{AB}, \sigma_{AB}) \leq F(\rho_A, \sigma_A) \quad \forall \rho_{AB}, \sigma_{AB}$ } HW

* **Joint concavity:** $F(\sum_i p_i \rho_i, \sum_i p_i \sigma_i) \geq \sum_i p_i F(\rho_i, \sigma_i)$ }

How about the trace distance? Monotonous with \geq , jointly convex via triangle ineq.

Fuchs-van de Graaf inequality: $1 - F \leq T \leq \sqrt{1 - F^2}$
 pure states } HW1

Channels in a classical world



How to describe mathematically? $y = f(x)$? In general, y can be random:

$$p(y|x) = \text{Pr}(\text{output } y \mid \text{input } x) \text{ s.t. } \begin{cases} p(y|x) \geq 0 \quad \forall x, y & \& \\ \sum_y p(y|x) = 1 \quad \forall x \end{cases}$$

* "Conditional prob distribution", "(memoryless) channel", "transition matrix", ...

* if input has distribution $p(x)$

$$\Rightarrow \text{output distribution } p(y) = \sum_x p(y|x) p(x)$$

} joint distribution $p(x, y) = p(x) p(y|x)$

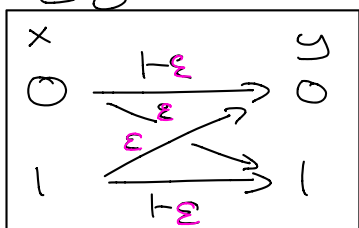
matrix-vector multiplication!

$$P_Y = P_{Y|X} \cdot P_X$$

$P_X \mapsto P_Y$ is linear!

Ex:

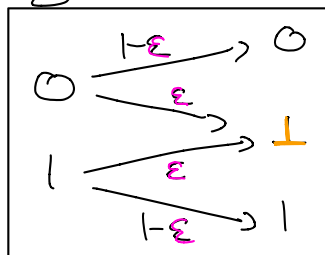
"binary symmetric channel"



$$\begin{aligned} p(0|0) &= 1 - \epsilon \\ p(1|1) &= 1 - \epsilon \\ p(1|0) &= \epsilon \\ p(0|1) &= \epsilon \end{aligned}$$

flip bit w/ probability ϵ

"binary erasure channel"



$$\begin{aligned} p(0|0) &= 1 - \epsilon \\ p(1|1) &= 1 - \epsilon \\ p(?|0) &= \epsilon \\ p(?|1) &= \epsilon \end{aligned}$$

bit lost w/ probability ϵ

\Rightarrow just like probability distributions can describe "uncertain states", channels can describe "uncertain" or "noisy" evolutions

GOAL of IT: Communicate reliably in the presence of uncertainty & noise!

How does the picture modify in quantum info theory?

Quantum Channels

$$D(\mathcal{H}_A) \ni \rho_A \rightarrow \boxed{?} \rightarrow \rho_B \in D(\mathcal{H}_B)$$

Since q. states are linear operators, this should be a map

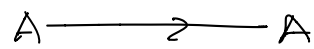
$$\Phi: L(\mathcal{H}_A) \rightarrow L(\mathcal{H}_B)$$

* Φ should be linear $\leadsto \Phi \in L(L(\mathcal{H}_A), L(\mathcal{H}_B))$, "superoperator"

Why? So that $\Phi[\sum_i p_i \rho_i] = \sum_i p_i \Phi[\rho_i]$

* NOTATION: $\Phi_{A \rightarrow B}[\rho_A]$ etc.

* Identity: $\mathcal{I}_A: L(\mathcal{H}_A) \rightarrow L(\mathcal{H}_A)$, $\mathcal{I}_A[\rho_A] = \rho_A$

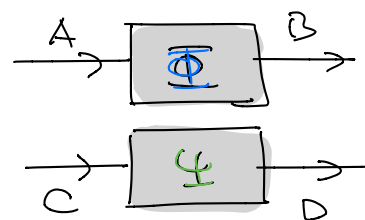


* Tensor product of superoperators:

$$L(\mathcal{H}_A \otimes \mathcal{H}_C) \rightarrow L(\mathcal{H}_B \otimes \mathcal{H}_D)$$

$$(\Phi_{A \rightarrow B} \otimes \Psi_{C \rightarrow D})[\rho_A \otimes \rho_C]$$

$$:= \Phi_{A \rightarrow B}[\rho_A] \otimes \Psi_{C \rightarrow D}[\rho_C]$$



& extend by linearity

What conditions should Φ satisfy to be a channel? (St attempt):

① Positivity: $\rho \geq 0 \Rightarrow \Phi[\rho] \geq 0 \quad \forall \rho$

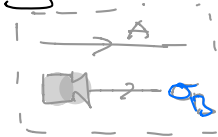
② Trace-preserving: $\text{tr}[\rho] = \text{tr}[\Phi[\rho]] \quad \forall \rho$

Together: ρ state $\Rightarrow \Phi[\rho]$ state

I write ρ since we will usually apply channels to states - but can plug in any linear operator!

Ex: * Base change: $\Phi[\rho] = U \rho U^\dagger$ for unitary (or isometry) U

* Add state: $\Phi[\rho_A] = \rho_A \otimes \sigma_B$ for state σ



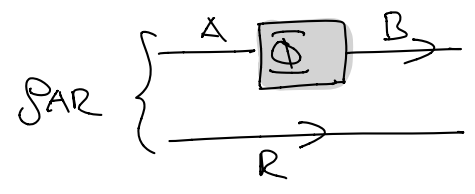
* Partial trace: $\Phi = \text{tr}_A$



* Measure-and-prepare: $\Phi[\rho] = \sum_x \langle x | \rho | x \rangle \cdot \sigma_x$



PROBLEM: There are superoperators s.t. ① hold for Φ , but NOT for some $\Phi \otimes \mathbb{I}_R$ ⚡



Transpose map: $\Upsilon[M] := M^T$ ← ^{w.r.t.} fixed basis

* Satisfies ① + ②

* $\rho_{AR} = |\Phi^+\rangle\langle\Phi^+| = \frac{1}{2} (|00\rangle\langle 00| + |11\rangle\langle 11| + |01\rangle\langle 10| + |10\rangle\langle 01|)$ Max ent state

$\Rightarrow (\Upsilon \otimes \mathbb{I}_R)[\rho_{AR}] = \frac{1}{2} (|00\rangle\langle 00| + |11\rangle\langle 11| + |10\rangle\langle 01| + |01\rangle\langle 10|)$

$= \frac{1}{2} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$ has eigenvalue -1 ⚡ $\begin{pmatrix} 0 \\ -1 \\ -1 \\ 0 \end{pmatrix} \leftrightarrow \begin{pmatrix} 0 \\ -1 \\ 1 \\ 0 \end{pmatrix}$

tr = 1

↳ ① is not good enough! The following def'n fixes this:

- Quantum channel: "superoperator" $\Phi_{A \rightarrow B} \in L(L(\mathcal{H}_A), L(\mathcal{H}_B))$ s.t.
- ① Completely positive: $\forall \mathcal{H}_R: M_{AR} \geq 0 \Rightarrow (\Phi_{A \rightarrow B} \otimes \mathbb{I}_R)[M_{AR}] \geq 0$
 - ② Trace-preserving: $\text{tr}[M_A] = \text{tr}[\Phi_{A \rightarrow B}[M_A]] \quad \forall M_A$

$\mathcal{C}(\mathcal{H}_A, \mathcal{H}_B) := \{\text{channels } \Phi_{A \rightarrow B}\}, \quad \mathcal{C}(\mathcal{H}_A) := \mathcal{C}(\mathcal{H}_A, \mathcal{H}_A)$

- * all examples above are channels — except the transpose → **EX & HW**
- * "dd from new": $\Phi_{A \rightarrow B} \otimes \Phi_{C \rightarrow D}$ & $\Phi_{B \rightarrow C} \circ \Phi_{A \rightarrow B}$ are again channels
- * any channel is composition of example channels ▽ → next week

NB: "complete positivity" not an issue for probability theory
 If $p(y|x)$ channel then so is $p(yz|xz') = p(y|x) \delta_{zz'}$

