

Holevo Bound & Relative Entropy

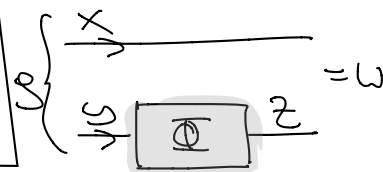
Last week: Entropy, Mutual info, properties

Recall: Strong subadditivity: $\forall \rho_{ABC} \in D(A \otimes B \otimes C)$:

$$I(A:B) \leq I(A:BC) \quad \text{i.e.} \quad H(AB) + H(BC) \geq H(B) + H(ABC)$$

Data Processing Ineq (DPI): $\forall \rho_{xy} \in D(\mathcal{X} \otimes \mathcal{Y}), \Phi \in C(\mathcal{Y}, \mathcal{Z})$:

$$I(X:Z)_\omega \leq I(X:Y)_\rho \quad \text{for } \omega_{xz} = (\mathcal{I}_x \otimes \Phi)[\rho_{xy}]$$



* generalizes SSA: $X=A, Y=BC, Z=B, \Phi = \text{tr}_C$

* implied by SSA: choose Stinespring isometry $V: \mathcal{Y} \rightarrow \mathcal{Z} \otimes \mathcal{W}$ for Φ :

$$\omega_{xzw} = (\mathcal{I}_x \otimes V) \rho_{xy} (\mathcal{I}_x \otimes V^\dagger) \text{ extends } \omega_{xz}$$

$$\& \quad I(X:Y)_\rho \stackrel{\text{Metric Invariance}}{=} I(X:ZW)_\omega \stackrel{\text{SSA}}{\geq} I(X:Z)_\omega \quad \square$$

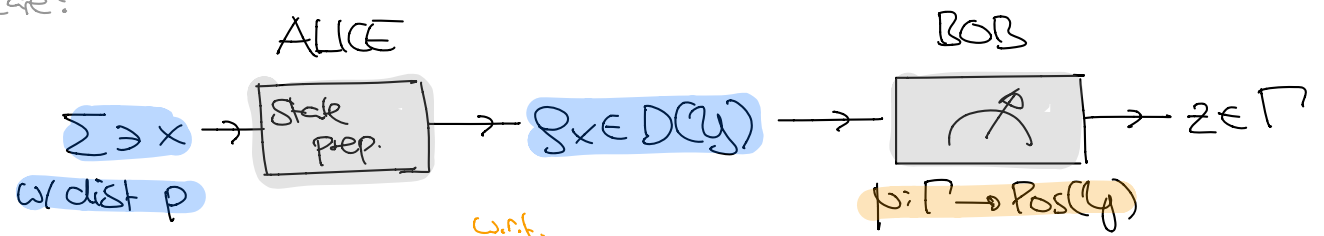
Recall: Ensembles $\{p_x, \rho_x\}$ can be represented by cq-states

$$\rho_{xy} = \sum_x p_x |x\rangle\langle x| \otimes \rho_x \in D(\mathcal{X} \otimes \mathcal{Y})$$

$$\text{Holevo } \chi\text{-quantity: } \chi(\{p_x, \rho_x\}) = I(X:Y) = H\left(\sum_x p_x \rho_x\right) - \sum_x p_x H(\rho_x)$$

* $\chi \leq H(X)$, with " $=$ " iff $\{\rho_x: p_x > 0\}$ pairwise orthogonal image \leadsto **ESET** !

Why care?



How large can $I(X:Z)$ be? $p(x,z) = p_x \cdot \text{tr}[\rho_x p(z)]$

Maximum over all p is known as "accessible information" of ensemble.

$$\text{Thm (Holevo): } I(X:Z) \leq \chi(\{p_x, \rho_x\}) = I(X:Y)$$

Pf: Apply DPI to ρ_{xy} and $\Phi[\sigma] = \sum_z \text{tr}[\sigma \rho(z)] |z\rangle\langle z|$ □

$$\left(\omega_{xz} = \sum_x p_x(x) |x\rangle\langle x| \otimes \Phi[\rho_x] = \sum_{x,z} p(x|z) |x\rangle\langle x| \otimes |z\rangle\langle z| \right)$$

Interpretation?

can decode X from Z
 (i.e. $X=f(Z)$)

ESET 7

$$\Leftrightarrow H(X|Z) = H(Z) \Leftrightarrow I(X:Z) = H(X)$$

"To communicate n random bits (perfectly), need to send n qubits."

$$\dim \mathcal{Y} \geq |\Sigma|$$

$p > 0$
 \Leftrightarrow
 $\{s_x\}$
 orthog.

$$\Downarrow \text{Holevo} \uparrow$$

$$X = H(X)$$

PSET

* quantitative bounds via Fano's inequality



Relative Entropy

Relative entropy: For $p, q \in \mathcal{P}(\Sigma)$, define

Consider w/ $\log(\dots + \epsilon)$ and $\epsilon \rightarrow 0$

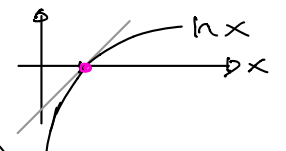
$$D(p||q) = \begin{cases} \sum_x p(x) \log \frac{p(x)}{q(x)} - \sum_x p(x) \log q(x) \\ \infty \end{cases}$$

if $\forall x: q(x) = 0 \Rightarrow p(x) = 0$
 otherwise

NB: $p(x) \cdot \log \frac{p(x)}{q(x)}$ clear if $p(x) > 0$, $= \begin{cases} 0 & \text{if } q(x) = p(x) = 0 \\ -\infty & \text{if } q(x) = 0, p(x) > 0 \end{cases}$

* $D(p||q) \geq 0$, $= 0$ iff $p=q$ \leftarrow asymmetric distance measure

Pf: wlog all $p(x) > 0$. Use $\ln(x) \leq x-1$, " $=$ " iff $x=1$.



$$D(p||q) = \sum_x p(x) \left(-\log \frac{q(x)}{p(x)} \right) \geq \frac{1}{\ln 2} \sum_x p(x) \left(1 - \frac{q(x)}{p(x)} \right) = 0$$

Why do we care?

"=" iff $q(x) = p(x) \forall x$ □

* $q(x) \equiv \frac{1}{|\Sigma|}$: $D(p||q) = \log |\Sigma| - H(p) \geq 0$, $= 0$ iff p uniform (8)

* $p_{xy} \in \mathcal{P}(\Sigma \times \Gamma)$, $q_{xy}(x,y) \equiv p_x(x) p_\gamma(y)$:

$$D(p_{xy}||q_{xy}) = \dots = I(X:Y)_p \geq 0, = 0 \text{ iff } p_{xy}(x,y) = p_x(x) p_\gamma(y) \quad (8)$$

Quantum relative entropy: $\rho, \sigma \in D(\mathcal{X})$

Consistent w/ $\log(-\text{tr} \cdot I)$ and $\epsilon \rightarrow 0$

$$D(\rho \parallel \sigma) = \begin{cases} \text{tr}[\rho \cdot \log \rho] - \text{tr}[\rho \cdot \log \sigma] & \text{if } \text{ker}(\sigma) \subseteq \text{ker}(\rho) \\ \infty & \text{otherwise} \end{cases}$$

equiv: $\text{im}(\rho) \subseteq \text{im}(\sigma)$

* $\rho \cdot \log(\sigma)$ clear on $\text{ker}(\sigma)^\perp$, define as zero on $\text{ker}(\sigma)$ if $\text{ker}(\sigma) \subseteq \text{ker}(\rho)$

↳ $D(\rho \parallel \sigma) < \infty$ iff $\text{ker}(\sigma) \subseteq \text{ker}(\rho)$ e.g. $D(|0\rangle\langle 0| \parallel (|+\rangle\langle +|)) = \infty$

* $\rho = \begin{pmatrix} p & \\ & \dots \end{pmatrix}, \sigma = \begin{pmatrix} q & \\ & \dots \end{pmatrix} : D(\rho \parallel \sigma) = D(p \parallel q)$

Monotonicity: $\rho, \sigma \in D(\mathcal{X}), \Phi \in C(\mathcal{X}, \mathcal{Y})$
 $\Rightarrow D(\rho \parallel \sigma) \geq D(\Phi[\rho] \parallel \Phi[\sigma])$

"FUNDAMENTAL THEOREM OF QIT"

PF: HARD ↙

Klein's inequality: $D(\rho \parallel \sigma) \geq 0, = 0$ iff $\rho = \sigma$

PF: Let $p: \Gamma \rightarrow \text{Pos}(\mathcal{X})$ measurement & $\Phi \in C(\mathcal{X}, \mathbb{C}^\Gamma)$ corresp. channel
 $\Phi[\omega] = \sum_{\gamma \in \Gamma} \text{tr}[\omega p(\gamma)] |\gamma\rangle\langle \gamma|$

$$\begin{cases} p(\gamma) = \text{tr}[\rho p(\gamma)] \\ q(\gamma) = \text{tr}[\sigma p(\gamma)] \end{cases} \quad \begin{cases} \Phi[\rho] = \sum_{\gamma} p(\gamma) |\gamma\rangle\langle \gamma| \\ \Phi[\sigma] = \sum_{\gamma} q(\gamma) |\gamma\rangle\langle \gamma| \end{cases}$$

a) $D(\rho \parallel \sigma) \geq D(\Phi[\rho] \parallel \Phi[\sigma]) = D(p \parallel q) \geq 0$

b) $\rho = \sigma : D(\rho \parallel \sigma) = 0 \rightarrow$ PSET

c) $\rho \neq \sigma : \exists p : \|p - q\|_1 \ominus \|\rho - \sigma\|_1 > 0 \Rightarrow D(\rho \parallel \sigma) \geq D(p \parallel q) \ominus 0 > 0 \quad \square$

↳ PSET / ESET for further properties

Applications: Quick proofs of things we proved last time

* $D(\rho \parallel \frac{I_X}{d}) \stackrel{\text{ESET}}{\geq} \log d - H(\rho) \sim H(\rho) \stackrel{\text{Klein}}{\leq} \log d, =$ iff $\rho = \frac{I_X}{d}$ ∞

* Subadditivity: $D(\rho_{XY} \parallel \rho_X \otimes \rho_Y) \stackrel{\text{ESET}}{\geq} I(X:Y)_{\rho_{XY}} \stackrel{\text{Klein}}{\geq} 0, = 0$ if $\rho_{XY} = \rho_X \otimes \rho_Y$ ∞

* Strong subadditivity:

$$I(X:YZ) = D(\rho_{XYZ} \parallel \rho_X \otimes \rho_{YZ}) \stackrel{\text{monotonicity}}{\geq} D(\rho_{XY} \parallel \rho_X \otimes \rho_Y) = I(X:Y) \quad \text{Q.E.D.}$$

$\Phi = \text{tr}_Z$

* Basis measurement: $\rho \in D(\mathbb{C}^{\mathcal{X}})$, $p(x) = \langle x | \rho | x \rangle$, $\sigma = \sum_x p(x) |x\rangle\langle x|$

$$\Rightarrow \boxed{H(\rho) = H(\sigma) \geq H(\rho), = \text{iff } \rho = \sigma}$$

Pf: $D(\rho \parallel \sigma) = -H(\rho) - \text{tr}[\rho \cdot \log(\sigma)] \stackrel{\text{Q.E.D.}}{=} -H(\rho) + H(\rho) \leadsto \text{KLEIN. } \square$

$\log(\sigma) = \sum_x \log(\langle x | \rho | x \rangle) |x\rangle\langle x|$