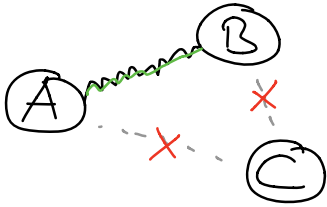


Monogamy of Entanglement



$$S_{AB} = |\psi_{AB}\rangle\langle\psi_{AB}|$$



$$S_{ABC} = S_{AB} \otimes S_C$$

$$\Rightarrow S_{AC} = S_A \otimes S_C \text{ (etc.)}$$

Can we make this more quantitative?

What is an entangled state?

Pure states: $|\psi_{AB}\rangle \neq |\psi_A\rangle \otimes |\psi_B\rangle \rightarrow$ Pset

$\Leftrightarrow S_A$ not pure

if state is pure

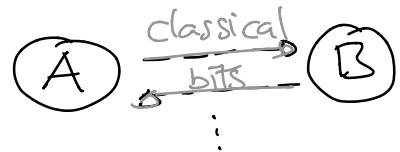
In general: S_{AB} is entangled if

$$S_{AB} \neq \sum_i p_i S_A^i \otimes S_B^i$$

MIXTURE;
not a superposition!

separable (unentangled) state

Why?



- All States that can be produced by LOCC (local ops & classical comm.)
- E.g., Classical probability distributions:

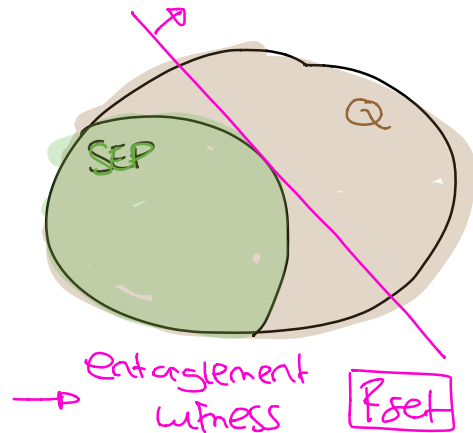
$$S_{AB} = \sum_{a,b} p(a,b) |a\rangle\langle a| \otimes |b\rangle\langle b|$$

• Convex geometry perspective:

$$\{\rho_{AB} \geq 0, \text{tr} \rho_{AB} = 1\} = \mathcal{Q}$$

$$\cup \{\rho_{AB} \text{ separable}\} =: \text{SEP}$$

Convex sets!



• NP-hard to test if given $\rho_{AB} \in \text{SEP}$!

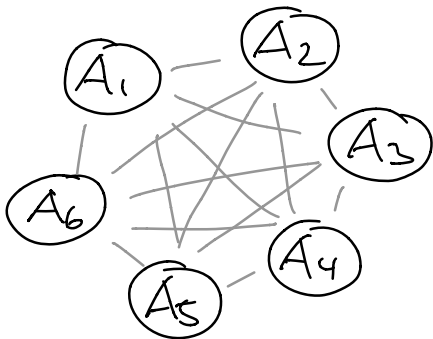
...but we will discuss the best known test...

Monogamy and Symmetry:

① In a permutation-symmetric state

$$|\psi_{A_1 \dots A_n}\rangle \in \text{Sym}^n(\mathbb{C}^d)$$

would expect that $S_{A_i A_j}$ "not very" entangled:



Ex: Mean-field $H = \sum_{i < j} (h_{ij})$
 $\hookrightarrow |E_0\rangle \in \text{Sym}^n$
 if nondegenerate

↑
Same interaction

Pset

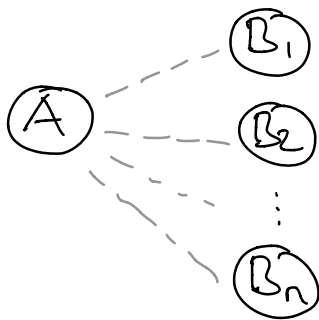
De Finetti: $S_{A_1 \dots A_k} \approx \int d\psi p(\psi) |\psi\rangle\langle\psi|^{\otimes k}$

E.g.:

$$|\Phi\rangle = \frac{1}{\sqrt{2}} (|000\rangle + |111\rangle) \text{ entangled!}$$

$$\Rightarrow \rho_{A_1 A_2} = \frac{1}{2} [|0\rangle\langle 0|^{\otimes 2} + |1\rangle\langle 1|^{\otimes 2}] \text{ NOT ENTANGLED !!!}$$

② If a state ρ_{AB_1} can be extended to $\rho_{AB_1 \dots B_n}$
 s.t. $\rho_{AB_1} = \dots = \rho_{AB_n}$ then ρ_{AB_1} "not very" entangled



$$\rho_{AB_1} \stackrel{?}{=} \sum_i p_i \rho_A^{(i)} \otimes \rho_{B_1}^{(i)}$$

distance?

Trace distance:

$$T(\rho, \sigma) = \max_{0 \leq Q \leq I} \text{tr}[Q(\rho - \sigma)] = \frac{1}{2} \|\rho - \sigma\|_1$$

• If $\rho = |\psi\rangle\langle\psi|$, $\sigma = |\phi\rangle\langle\phi|$ pure:

$$T(\rho, \sigma) = \sqrt{1 - |\langle\psi|\phi\rangle|^2}$$

Post 1

overlap!

trace norm:

$$\|\Delta\|_1 = \sum_i |\lambda_i|$$

$$\text{for } \Delta = \sum_i \lambda_i |e_i\rangle\langle e_i|$$

① Proof of the de Finetti theorem

$$|\Phi\rangle \in \text{Sym}^{k+n}(\mathbb{C}^d) \quad \text{MORE CONVENIENT}$$

Idea: Measure $\{Q_\psi\}$ on last n systems

→ first k should be $|\psi\rangle^{\otimes k}$

$$\begin{aligned} \underline{\rho}_k &= \text{tr}_n \left[|\Phi\rangle\langle\Phi| \right] \stackrel{\downarrow}{=} \text{tr}_n \left[(\mathbb{I}_k \otimes \Pi_n) |\Phi\rangle\langle\Phi| \right] \\ &= \binom{n+d-1}{n} \int d\psi \text{tr}_n \left[(\mathbb{I}_k \otimes |\psi\rangle^{\otimes n} \langle\psi|^{\otimes n}) |\Phi\rangle\langle\Phi| \right] \\ &\stackrel{\text{cyclic}}{=} \binom{n+d-1}{n} \int d\psi (\mathbb{I}_k \otimes \langle\psi|^{\otimes n}) |\Phi\rangle\langle\Phi| (\mathbb{I}_k \otimes |\psi\rangle^{\otimes n}) \\ &= \int d\psi p(\psi) |\psi\rangle\langle\psi| \end{aligned}$$

where we introduced

$$\sqrt{p(\psi)} \cdot |\psi\rangle = \binom{n+d-1}{n}^{1/2} (\mathbb{I}_k \otimes \langle\psi|^{\otimes n}) |\Phi\rangle$$

↑ prob. density ↑ unit vector

Claim: $\rho_k = \int d\psi p(\psi) |\psi\rangle\langle\psi| \approx \int d\psi p(\psi) |\psi\rangle^{\otimes k} \langle\psi|^{\otimes k}$

Average overlap: Will justify later!

$$\begin{aligned}
 & \int d\psi p(\psi) |\langle V_\psi | \psi^{\otimes k} \rangle|^2 \\
 &= \int d\psi p(\psi) \langle V_\psi | \psi^{\otimes k} \rangle \langle \psi^{\otimes k} | V_\psi \rangle \\
 &= \binom{n+d-1}{n} \int d\psi |\langle \Phi | \psi^{\otimes (n+k)} \rangle|^2 \\
 &= \binom{n+d-1}{n} \binom{n+k+d-1}{n+k}^{-1} \underbrace{\langle \Phi | \Pi_{n+k} | \Phi \rangle}_{=1}
 \end{aligned}$$

LAST

$$\geq 1 - \frac{kd}{n}$$

THU

Almost done: Could be on the PSET

$$\begin{aligned}
 & T(S_k, \int d\psi p(\psi) |\psi^{\otimes k} \rangle \langle \psi^{\otimes k}|) \\
 &= \frac{1}{2} \left\| \int d\psi p(\psi) [|V_\psi\rangle \langle V_\psi| - |\psi^{\otimes k}\rangle \langle \psi^{\otimes k}|] \right\|_1 \\
 &\leq \int d\psi p(\psi) \frac{1}{2} \left\| |V_\psi\rangle \langle V_\psi| - |\psi^{\otimes k}\rangle \langle \psi^{\otimes k}| \right\|_1 \\
 &= \int d\psi p(\psi) \sqrt{1 - |\langle V_\psi | \psi^{\otimes k} \rangle|^2} \\
 &\stackrel{\text{JENSEN}}{\leq} \left[\int d\psi p(\psi) (1 - |\langle V_\psi | \psi^{\otimes k} \rangle|^2) \right]^{1/2}
 \end{aligned}$$


$$= \left[1 - \underbrace{\int d\psi p(\psi) |\langle V_\psi | \psi^{\otimes k} \rangle|^2}_{\text{above!!!}} \right]^{1/2}$$

$$\leq \left[1 - \left(1 - \frac{kd}{n} \right) \right]^{1/2} = \sqrt{\frac{kd}{n}}$$

$$\Rightarrow \rho_k \approx \int d\psi p(\psi) |\psi\rangle\langle\psi|^{\otimes k} \quad \text{if } k \ll n \quad \square$$

Aside: De Finetti is useful in Q. crypto!

• general attacks \rightarrow "collective" attacks

• DI-QKD:  Bell tests at random sites

Same if ρ is mixed state on $\text{Sym}^N(\mathbb{C}^d)$. ($N=nk$)

What if only $[\mathbb{R}_\pi, \rho] = 0$? \leftarrow "Permutation-invariant"

PSET?

Fact: Exists purification $|\Phi\rangle \in \text{Sym}^N(\mathcal{H}_A \otimes \mathcal{H}_A)$!

$$\hookrightarrow \Phi_k \approx \int d\psi_{AA'} p(\psi) |\psi_{AA'}\rangle^{\otimes k} \langle\psi_{AA'}|^{\otimes k}$$

$$\hookrightarrow \rho_k \approx \int d\rho_{SA} p(\rho) \rho^{\otimes k} \quad \text{up to } \frac{kd^2}{n}$$

② Extendibility

If S_{AB} separable:

$$S_{AB} = \sum_i p_i S_A^i \otimes S_B^i$$

$$\leadsto S_{AB, \dots, B_n} = \sum_i p_i S_A^i \otimes (S_B^i)^{\otimes n}$$

"Symmetric" n-extendibility

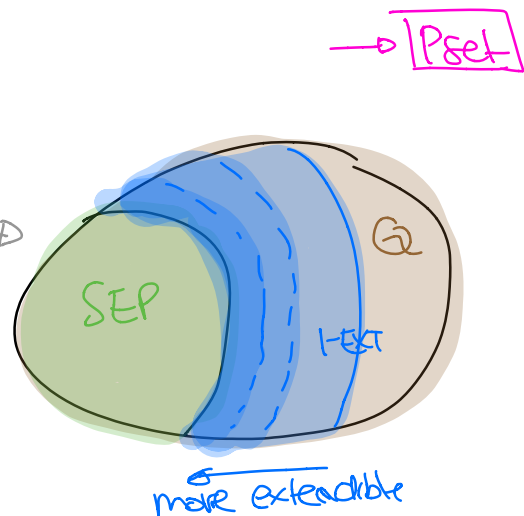
$$[I_A \otimes R_{\pi, S}] = 0$$

IF NOT SYMM: SYMMETRIZE!

Claim: If S_{AB} is symm. n -extendible then $O(\frac{1}{n})$ close to SEP!

• GOOD ENTANGLEMENT TEST!

↳ Semidefinite program



• Optimal? \rightarrow PSET