

# Message Passing for Decoding and Inference (§21/25/26)

## ① Decoding Problem:

$$s \rightarrow x^N \rightarrow y^N \xrightarrow{??} \hat{s}$$

Given channel  $Q$ , encoder, and output  $y^N$ , how to decode?

\* Clever algebra, like for Reed-Solomon: Not always possible!

\* Maximum likelihood codeword decoder: Assume  $P(s)$  uniform.

Given  $y^N$ , find  $s$  that maximizes  $P(s|y^N)$ .

Eqv: Maximize  $P(x^N|y^N)$  over  $C = \{x^N \text{ codeword}\}$  We assume no two  $s$  have same codeword  $x^N$  !!!

$$P(x^N|y^N) \stackrel{\text{Bayes}}{=} \frac{P(y^N|x^N) P(x^N)}{P(y^N)} = \frac{1}{P(y^N) \cdot \#C} P(y^N|x^N) \cdot \mathbb{1}[x^N \in C]$$

Can ignore if  $y^N$  known  
remember this notation?

Thus: Given  $y^N$ , find  $x^N$  that maximizes

$$G(x^N) = Q(y_1|x_1) \dots Q(y_N|x_N) \mathbb{1}[x^N \in C]$$

e.g. repetition code  $R_3$ :  $C = \{000, 111\}$

$$G(x_1, x_2, x_3) = Q(y_1|x_1) Q(y_2|x_2) Q(y_3|x_3) \delta_{x_1, x_2} \delta_{x_2, x_3}$$

notation ok?

\* Bitwise decoder: Maximize  $P(x_i|y^N)$  for each  $i=1..N$ . Equivalently:

Given  $y^N$ , find  $x_i$  that maximizes

$$G_i(x_i) := \sum_{\{x^N \mid x_j \neq i\}} G(x^N)$$

How to avoid having to compute all numbers  $G(x^N)$  ???

Exponential in  $N$

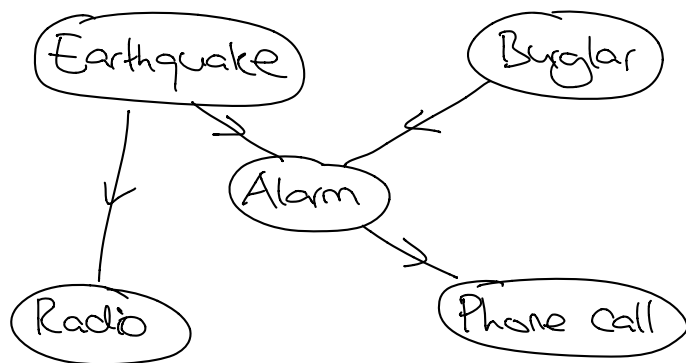
## ② Inference in Bayesian networks:

Assume we have a model

$$P(e, b, r, a, p)$$

$$= \underbrace{P(e)} \underbrace{P(b)} \underbrace{P(r|e)} \underbrace{P(a|e, b)} \underbrace{P(p|a)}$$

all functions are known (= model)



## Inference Problems:

$$? = \Pr(B=1 | A=1) = \frac{P_r(B=1, A=1)}{\sum_b P_r(B=b, A=1)}$$

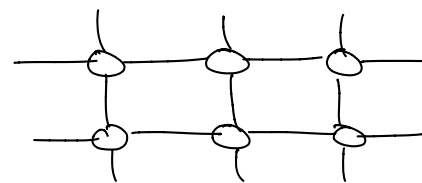
↳ enough to compute  $\underbrace{P_r(B=b, A=1)}_{\text{marginal of } B} = \sum_{e, r, p} \underbrace{P(e, b, r, p)}_{\text{product of local factors}}$

How to avoid having to first compute  $P(e, b, r, p)$  for all  $e, b, r, p$ ?

## ③ Statistical Physics: Ising model on lattice

\* one particle per site with states  $x_i \in \{\pm 1\}$

\* total energy  $E[\{x_i\}] = \sum_{i \sim j} \frac{J}{2} (1 - x_i x_j)$



energy cost  $J$  if NOT same  
ferromagnet if  $J > 0$

Partition function at temperature  $T$ :

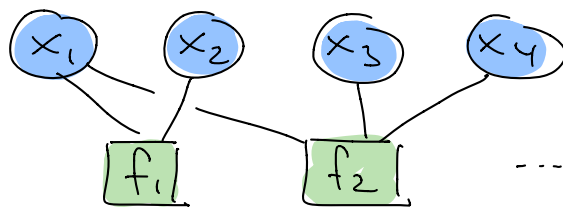
$$Z = \sum_{\{x_i\}} e^{-E[\{x_i\}]/T} = \sum_{\{x_i\}} \underbrace{\prod_{i \sim j} e^{-J/2T \cdot (1 - x_i x_j)}}_{\text{product of local factors}}$$

# General Setup:

$$G(x^N) = \prod_m \text{factor } f_m(\{x_i\}_{i \in I(m)})$$

Subset of variables

Where  $x_i \in \mathcal{X}_i$



$$G(x^N) = f_1(x_1, x_2) f_2(x_1, x_3, x_4) \dots$$

## Factor graph:

\* vertices:  $x_i$  for each variable,  $f_m$  for each factor

\* edge:  $x_i - f_m$  if  $f_m$  depends on  $x_i$

## Problem:

Compute marginals

→ ① Bitwise decoding

② Bayesian inference...

$$G_i(x_i) := \sum_{\{x_j\}_{j \neq i}} G(x^N) = \sum_{\substack{x_1 \dots x_{i-1} \\ x_{i+1} \dots x_N}} G(x_1, \dots, x_N)$$

e.g. for repetition code: If we receive  $y^3 = 110$ :

$$G(x_1, x_2, x_3) = \delta_{x_1, x_2} \delta_{x_2, x_3} Q(1|x_1) Q(1|x_2) Q(0|x_3)$$

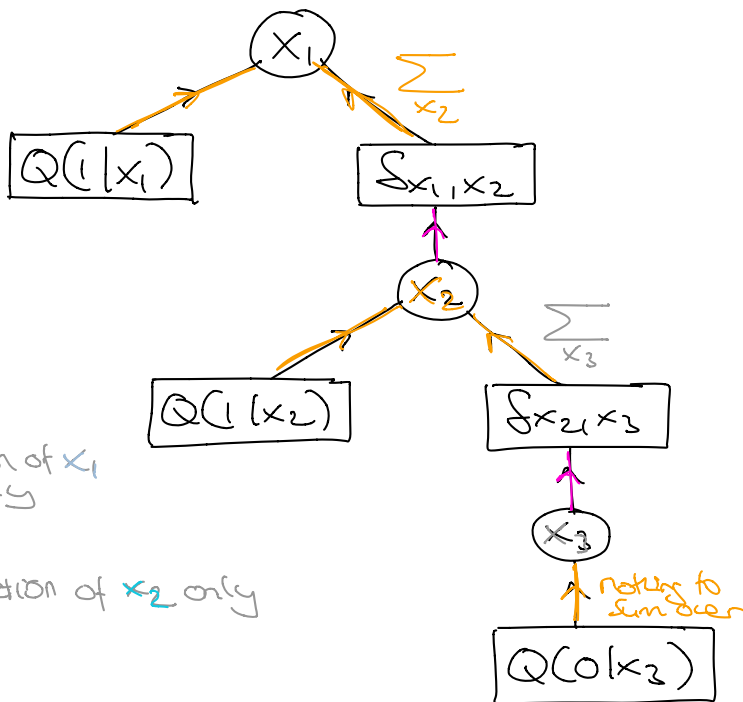
\* factor graph is tree:

\* if rooted at  $x_1$ , gives natural "algo" for computing marginal:

$$G_1(x_1) = Q(1|x_1)$$

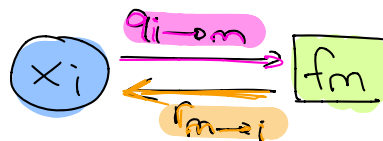
$$\sum_{x_2} \delta_{x_1, x_2} Q(1|x_2) \left. \vphantom{\sum_{x_2}} \right\} \text{function of } x_1 \text{ only}$$

$$\sum_{x_3} \delta_{x_2, x_3} Q(0|x_3) \left. \vphantom{\sum_{x_3}} \right\} \text{function of } x_2 \text{ only}$$



⚡ If interested in  $G_2$  or  $G_3$ , need to change root and run again

The following "message passing" algorithm computes all  $G_i(x_i)$  at the same time:



# Sum-product algorithm ("belief propagation"):

Input: Factor graph & factors  $\{f_i\}$  & integer  $T$

① For all edges  $(x_i) - [f_m]$  and all  $x_i \in \mathcal{X}_i$ :

$$q_{i \rightarrow m}(x_i) \leftarrow 1$$

② For  $T$  steps:

For all edges  $(x_i) - [f_m]$  and all  $x_i \in \mathcal{X}_i$ :

$$r_{m \rightarrow i}(x_i) \leftarrow \sum_{\{x_j\}_{j \in I(m), j \neq i}} f_m(\{x_j\}_{j \in I(m)}) \cdot \prod_{j \in I(m), j \neq i} q_{j \rightarrow m}(x_j)$$

$I(m)$  = variables that appear in  $f_m$

For all edges  $(x_i) - [f_m]$  and all  $x_i \in \mathcal{X}_i$ :

$$q_{i \rightarrow m}(x_i) \leftarrow \prod_{n \in \mathcal{N}(i), n \neq m} r_{n \rightarrow i}(x_i)$$

$\mathcal{N}(i)$  = factors that depend on  $x_i$

③ For all vertices  $(x_i)$  and all  $x_i \in \mathcal{X}_i$ :

$$G_i(x_i) \leftarrow \prod_{m \in \mathcal{N}(i)} r_{m \rightarrow i}(x_i)$$

NB: Messages are functions/tuples! (one real number for each  $x_i \in \mathcal{X}_i$ )

\* Sum-product algo works provably for trees

\* computes all  $G_i$  at the same time if  $T \geq$  diameter of graph.

\* in practice also used for general graphs

but only a heuristic: problem is **NP-hard**

## Variations:

\* Partition function:  $Z = \sum_{x^N} G(x^N) = ? \rightarrow$  ③ physics  $[Z = \sum_{x_i} G_i(x_i)]$

\* Maximum:  $\max_{x^N} G(x^N) = ? \rightarrow$  ① tic decoding

[Replace  $\sum$  by  $\max \rightsquigarrow$  "max-product" algo  $\xrightarrow{-\log}$  "min-sum" algo]

## Outlook

What did we NOT cover?

- \* Channels with memory
- \* Multi-user information theory
- \* More connections to inference, machine learning, etc.
- \* Quantum information theory → MasterMath course with Mari's Ozols & Quantum Computing → courses by Mari (BSc) + Ronald de Wolf (MSc)
- \* Connections to Cryptography → course by Chris Schaffner

Bachelor project?

## How to prepare for the exam?

- \* Learning objectives @ homepage
- \* Lecture notes, homework, practice problems, last year's exam
- \* Don't forget to prepare your cheat sheet
- \* Structure: mix of problems of type ① + ② from HW

THANKS + SEE YOU AGAIN SOON !