

# Introduction to Information Theory, Fall 2020

## Practice problems for exercise class #4

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You do **not** have to hand in these exercises, they are for your practice only.

0. Exercises from **MacKay**: 5.21 (only for  $X^2$  and  $X^3$ ), 5.25

1. **Relative entropy** Given two probability distributions  $P$  and  $Q$  on the same set  $\mathcal{A}_X$ , we define the *relative entropy*  $D(P \parallel Q)$  by

$$D(P \parallel Q) = \sum_{x \in \mathcal{A}_X} P(x) \log\left(\frac{P(x)}{Q(x)}\right). \quad (1)$$

- (a) In what situation is  $D(P \parallel Q)$  infinite?
- (b) Show that the (ordinary) entropy  $H(P)$  can be written as

$$H(P) = \log|\mathcal{A}_X| - D(P \parallel U)$$

where  $U$  is the uniform distribution over  $\mathcal{A}_X$ .

- (c) Show that  $D(P \parallel Q) \geq 0$ .  
*Hint: Note that this is the same as the Gibbs inequality that we proved in class. Try proving it again to practice using Jensen's inequality.*
- (d) Show that  $D(P \parallel Q) = 0$  if and only if  $P = Q$  (and carefully distinguish the case where  $P(x) = 0$  for some  $x \in \mathcal{A}_X$ ).

2. **Using the wrong symbol code** Suppose we are given two probability distributions  $P$  and  $Q$  on the same set  $\mathcal{A}_X$ . We know that there exists a uniquely decodable symbol code with codeword lengths  $l(C(x)) = \lceil \log \frac{1}{Q(x)} \rceil$  and it has expected length satisfying

$$H(Q) \leq L(C, Q) < H(Q) + 1.$$

Show that if we use this code for the 'wrong' distribution  $P$  the expected codelength will satisfy

$$H(P) + D(P \parallel Q) \leq L(C, P) < H(P) + D(P \parallel Q) + 1$$

where the relative entropy is defined by Eq. (1).

3. **Optimality of the Huffman code (mathematics challenge):**

- (a) For any probability distribution show that there exists an optimal code  $C$  with the following properties:
  - If  $p(x) > p(y)$  then  $l(C(x)) \leq l(C(y))$
  - The two longest codewords have the same length.
  - Two of the longest codewords differ only in the last bit and correspond to two of the least likely symbols.
- (b) Show that the Huffman code is optimal. (If you want, you can look up the solution to this exercise in the book 'Elements of Information Theory' by Cover and Thomas.)