

Message Passing for Decoding and Inference (§21/25/26)

Recall from Tue/Thu:

① Decoding Problem:

$$s \rightarrow x^N \rightarrow y^N \xrightarrow{???} \hat{s}$$

Given channel Q , encoder, and output y^N , how to decode?

* Maximum likelihood codeword decoder: Assume $P(s)$ uniform. Given y^N , find s that maximizes $P(s|y^N)$. Equ: Maximize $P(x^N|y^N)$ over $C = \{x^N \text{ codeword}\}$. We assume **no two s have the same codeword x^N !!!**

$$P(x^N|y^N) \stackrel{\text{Bayes}}{=} \frac{P(y^N|x^N) P(x^N)}{P(y^N)} = \frac{1}{P(y^N) \cdot \#C} P(y^N|x^N) \cdot \mathbb{1}[x^N \in C]$$

Can ignore if y^N known
remember this notation?

Thus: Given y^N , find x^N that maximizes

$$G(x^N) = Q(y_1|x_1) \dots Q(y_N|x_N) \mathbb{1}[x^N \in C]$$

* Bitwise decoder: Maximize $P(x_i|y^N)$ for each $i=1, \dots, N$. Equivalently: Given y^N , find x_i that maximizes

$$G_i(x_i) := \sum_{\{x_j\}_{j \neq i}} G(x^N)$$

How to avoid having to compute all numbers $G(x^N)$???

Exponential in N

e.g. repetition code R_3 : $C = \{000, 111\}$

$$G(x_1, x_2, x_3) = Q(y_1|x_1) Q(y_2|x_2) Q(y_3|x_3) \underbrace{\delta_{x_1, x_2} \delta_{x_2, x_3}}_{\text{notation ok?}}$$

e.g. linear code with parity check matrix H : $C = \{x^N : Hx^N = 0\}$

$$G(x^N) = Q(y_1|x_1) \dots Q(y_N|x_N) \mathbb{1}[Hx^N = 0]$$

or, in terms of noise vector $n^N = x^N \oplus y^N$:

$$\tilde{G}(n^N) = P(n_1) \dots P(n_N) \cdot \mathbb{1}[Hn^N = Hy^N]$$

assuming $y^N = x^N \oplus n^N$ $\xrightarrow{\text{i.i.d. noise}}$

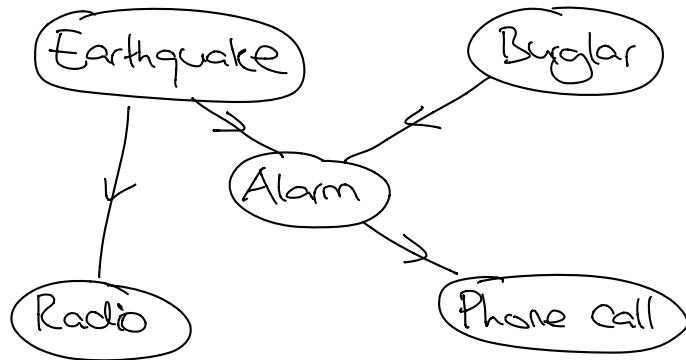
② Inference in Bayesian networks:

Assume we have a model

$$P(e, b, r, a, p)$$

$$= \underbrace{P(e)} \underbrace{P(b)} \underbrace{P(r|e)} \underbrace{P(a|e, b)} \underbrace{P(p|a)}$$

all functions are known (= model)



Inference Problems:

$$? = \Pr(B=1 | A=1) = \frac{\Pr(B=1, A=1)}{\sum_b \Pr(B=b, A=1)}$$

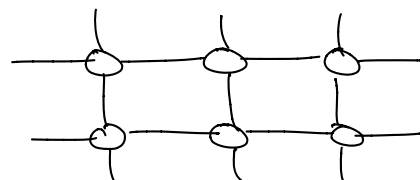
↳ enough to compute $\underbrace{\Pr(B=b, A=1)}_{\text{marginal of } B} = \sum_{e, r, p} \underbrace{P(e, b, r, 1, p)}_{\text{product of local factors}}$

How to avoid having to first compute $P(e, b, r, 1, p)$ for all e, b, r, p ?

③ Statistical Physics: Ising model on lattice

* one particle per site with states $x_i \in \{\pm 1\}$

* total energy $E[\{x_i\}] = \sum_{i \sim j} \frac{J}{2} (1 - x_i x_j)$
 energy cost J if NOT same
 ferromagnet if $J > 0$



Partition function at temperature T :

$$Z = \sum_{\{x_i\}} e^{-E[\{x_i\}]/T} = \sum_{\{x_i\}} \prod_{i \sim j} e^{-J/2T \cdot (1 - x_i x_j)}$$

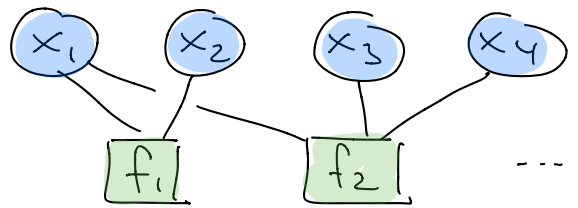
product of local factors

General Setup:

$$G(x^N) = \prod_m \overset{\text{factor}}{f_m(\{x_i\}_{i \in I(m)})}$$

Subset of variables

Where $x_i \in \mathcal{X}_i$



$$G(x^N) = f_1(x_1, x_2) f_2(x_1, x_3, x_4) \dots$$

Factor graph:

* vertices: x_i for each variable, f_m for each factor

* edge: $x_i - f_m$ if f_m depends on x_i

Problem:

Compute marginals

① Bitwise decoding, Bayesian inference, ... ②

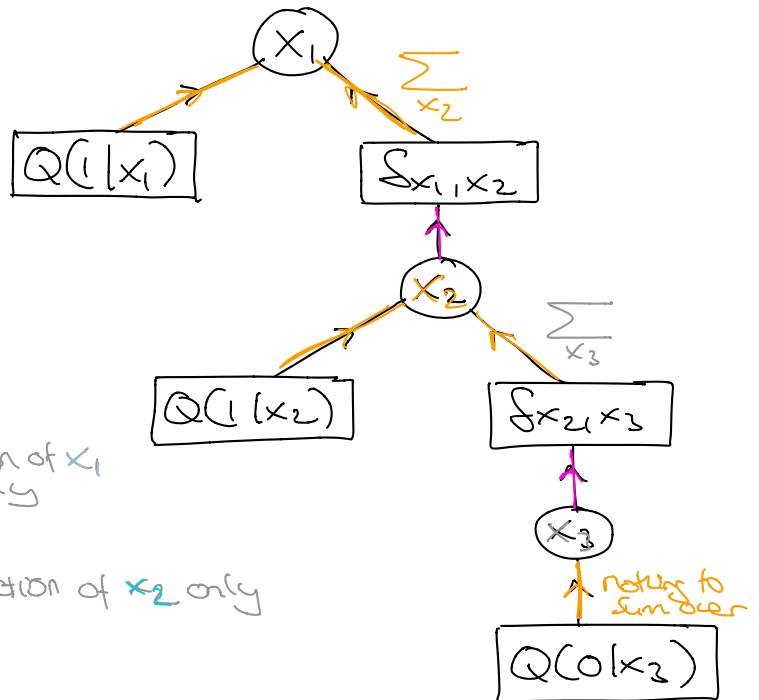
$$G_i(x_i) := \sum_{\{x_j\}_{j \neq i}} G(x^N) = \sum_{\substack{x_1, \dots, x_{i-1} \\ x_{i+1}, \dots, x_N}} G(x_1, \dots, x_N)$$

e.g. for repetition code: If we receive $y^3 = 110$:

$$G(x_1, x_2, x_3) = \delta_{x_1, x_2} \delta_{x_2, x_3} Q(1|x_1) Q(1|x_2) Q(0|x_3)$$

* factor graph is tree:

* if rooted at x_1 , gives natural "algo" for computing marginal:



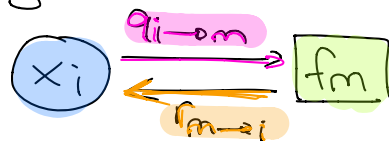
$$G_1(x_1) = Q(1|x_1)$$

$$\sum_{x_2} \delta_{x_1, x_2} Q(1|x_2) \left. \vphantom{\sum_{x_2}} \right\} \text{function of } x_1 \text{ only}$$

$$\sum_{x_3} \delta_{x_2, x_3} Q(0|x_3) \left. \vphantom{\sum_{x_3}} \right\} \text{function of } x_2 \text{ only}$$

⚡ If interested in G_2 or G_3 , need to change root and run again

↳ "decentral" message passing algo



Sum-product algorithm ("belief propagation"):

Input: Factor graph & factors $\{f_i\}$ & integer T

① For all edges $(x_i) - [f_m]$ and all $x_i \in \mathcal{X}_i$:

$$q_{i \rightarrow m}(x_i) \leftarrow 1$$

② For T steps:

For all edges $(x_i) - [f_m]$ and all $x_i \in \mathcal{X}_i$:

$$r_{m \rightarrow i}(x_i) \leftarrow \sum_{\{x_j\}_{j \in I(m), j \neq i}} f_m(\{x_j\}_{j \in I(m)}) \cdot \prod_{j \in I(m), j \neq i} q_{j \rightarrow m}(x_j)$$

$I(m)$ = variables that appear in f_m

For all edges $(x_i) - [f_m]$ and all $x_i \in \mathcal{X}_i$:

$$q_{i \rightarrow m}(x_i) \leftarrow \prod_{n \in \Pi(i), n \neq m} r_{n \rightarrow i}(x_i)$$

$\Pi(i)$ = factors that depend on x_i

③ For all vertices (x_i) and all $x_i \in \mathcal{X}_i$:

$$G_i(x_i) \leftarrow \prod_{m \in \Pi(i)} r_{m \rightarrow i}(x_i)$$

NB: Messages are functions/tuples! (one real number for each $x_i \in \mathcal{X}_i$)

* Sum-product algo works provably for trees

* computes all G_i at the same time if $T \geq$ diameter of graph.

* in practice also used for general graphs \rightarrow Lecture 13

but only a heuristic: problem is **NP-hard**

Variations:

* Partition function: $Z = \sum_{x^N} G(x^N) = ? \rightarrow$ ③ physics $[Z = \sum_{x_i} G_i(x_i)]$

* Maximum: $\max_{x^N} G(x^N) = ? \rightarrow$ ① tic decoding

[Replace \sum by $\max \rightsquigarrow$ "max-product" algo $\xrightarrow{-\log}$ "min-sum" algo]

Outlook

What did we NOT cover?

- * Channels with memory
 - * Multi-user information theory
 - * More connections to inference, machine learning, etc.
 - * How to decode barcode from Lecture 1 ☹️
 - * Quantum information theory → MasterMath course with Maris Ozols & Quantum Computing → courses by Maris (BSc) + Ronald de Wolf (MSc)
 - * Connections to Cryptography → course by Chris Schaffner
- } Baddeley project?

How to prepare for the exam?

- * Learning objectives @ homepage course manual NOT up to date!!!
- * Lecture notes, homework, problems from the ex. class
- * Don't forget to prepare your cheat sheet
- * Structure: mix of problems of type ① + ② from HW

THANKS + SEE YOU AGAIN SOON! ⚠️