- 1. (4 points) Fidelity inequality: Let  $\mathcal{X}$  be a complex Euclidean space with  $\dim(\mathcal{X}) \geq 2$ .
  - (a) Let  $|u_1\rangle, |u_2\rangle, |v\rangle \in \mathcal{X}$  be arbitrary pure quantum states. Show that

$$|\langle u_1|v\rangle|^2 + |\langle u_2|v\rangle|^2 \le 1 + |\langle u_1|u_2\rangle|.$$

Hint: Upper bound the left-hand side by the largest eigenvalue of some rank-2 matrix. Compute this eigenvalue to get the right-hand side.

(b) Let  $\rho_1, \rho_2, \sigma \in D(\mathcal{X})$  be arbitrary states. Show that

$$F(\rho_1, \sigma)^2 + F(\rho_2, \sigma)^2 \le 1 + F(\rho_1, \rho_2).$$

- 2. (4 points) Entanglement cost using compression and teleportation: In this exercise you will give an alternative proof for the fact that the entanglement cost is at most the entanglement entropy for a pure state. Let  $|\psi\rangle_{XY} \in \mathcal{S}(\mathcal{X} \otimes \mathcal{Y})$  be a pure state.
  - (a) Let  $\rho_X$  and  $\rho_Y$  be the reduced density matrices and let  $\alpha > H(\rho_X) = H(\rho_Y)$ . Show, using compression, that for all  $\delta > 0$  there exists an LOCC protocol for all but finitely many n which converts  $\lfloor \alpha n \rfloor$  Bell pairs into a state  $\tilde{\psi}_n$  with  $F(\psi^{\otimes n}, \tilde{\psi}_n) > 1 \delta$ .

    Hint: Use teleportation!
  - (b) Use (a) to show that for every pure state  $E_C(X : Y) \leq H(X)$ .
- 3. (3 points) Entanglement rank and the fidelity with the maximally entangled state: Let  $\Sigma = \{1, ..., n\}$  and  $\mathcal{X} = \mathcal{Y} = \mathbb{C}^{\Sigma}$  be two complex Euclidean spaces of dimension n. Let  $\tau_n \in D(\mathcal{X} \otimes \mathcal{Y})$  be the canonical maximally entangled state, i.e.,  $\tau_n = |\tau_n\rangle\langle\tau_n|$  where

$$|\tau_n\rangle = \frac{1}{\sqrt{n}} \sum_{i=1}^n |i\rangle \otimes |i\rangle.$$

Show that, for any state  $\rho \in D(\mathcal{X} \otimes \mathcal{Y})$  of entanglement rank r,  $F(\tau_n, \rho)^2 \leq r/n$ .

4. (5 points) Entanglement distillation: For any  $p \in [0, 1]$ , let

$$|\tau(p)\rangle_{\mathsf{AB}} = \sqrt{p}\;|0\rangle\otimes|0\rangle + \sqrt{1-p}\;|1\rangle\otimes|1\rangle$$

be a two-qubit state shared between Alice and Bob, and let  $|\tau\rangle_{AB} = |\tau(1/2)\rangle_{AB}$  denote the maximally entangled state. In the following, we always assume that Alice and Bob are given  $|\tau(p)\rangle_{AB}^{\otimes n}$ , for some  $n \geq 1$  and  $0 \leq p \leq 1/2$ , as their input.

(a) Using LOCC, they would like to distill one perfect copy of  $|\tau(q)\rangle_{AB}$  that is as close as possible to the maximally entangled state  $|\tau\rangle_{AB}$ . Derive a formula for the largest possible  $q \leq 1/2$  they can get and express it as an explicit function of p and n. What is the maximal fidelity their output state  $|\tau(q)\rangle_{AB}$  can have with the desired target state  $|\tau\rangle_{AB}$ , i.e., what is

$$\max_{q} F(|\tau(q)\rangle, |\tau\rangle),$$

where the maximum is over all possible achievable q?

- (b)  $\blacksquare$  Take p = 0.1 and plot this as a function of n. Determine the smallest n such that  $|\tau(p)\rangle_{\mathsf{AB}}^{\otimes n}$  can be perfectly transformed into  $|\tau\rangle_{\mathsf{AB}}$  by LOCC.
- (c) This time Alice and Bob want to extract more entanglement. That is, they want to perfectly obtain the state

$$|\tau_d\rangle_{\mathsf{A}'\mathsf{B}'} = \frac{1}{\sqrt{d}} \sum_{i=1}^d |i\rangle_{\mathsf{A}'} |i\rangle_{\mathsf{B}'},$$

for some  $d \geq 2$ . Derive a formula for the largest possible d as a function of p and n.

- (d)  $\blacksquare$  Take p = 0.1 and plot this as a function of n. Determine the smallest n such that  $|\tau(p)\rangle_{\mathsf{AB}}^{\otimes n}$  can be perfectly transformed into three copies of  $|\tau\rangle_{\mathsf{AB}}$  by LOCC.
- (e)  $\blacksquare$  Take p = 0.2 and n = 2. What is the largest possible fidelity their output state can have with  $|\tau_3\rangle_{AB}$ ?