

Last time: Quantum Schur transform

key ingredient, still need to construct!

A circuit for the Clebsch-Gordan transform

$$J_k: \text{Sym}^k(\mathbb{C}^2) \otimes \mathbb{C}^2 \stackrel{\cong}{=} \bigoplus_{p=\pm 1} \text{Sym}^{k+p}(\mathbb{C}^2)$$

CG coeffs  $\nearrow$

How to construct? Recall how we proved  $\cong \dots$

key tool:  $\{R_u\} \rightsquigarrow r_z := -i\partial_s = d_{R_{e^{is}z}}$

\* Why?  $J_k r_z^{\text{LHS}} J_k^\dagger = r_z^{\text{RHS}} \rightarrow$  preserves eigenspaces

So let's diagonalize  $r_z^{\text{LHS}}$  &  $r_z^{\text{RHS}}$  ...

\*  $\text{Sym}^k(\mathbb{C}^2)$ :  $T_u^{(k)} \rightsquigarrow t_z^{(k)} = \text{restr. of } z \otimes I \otimes \dots \otimes I + \dots + I \otimes \dots \otimes I \otimes z$

NOTATION:  $|k, s\rangle := |w_{m, k-m}\rangle$  eigenvectors

$s := 2m - k \in \{k, k-2, \dots, -k\}$  eigenvalues

discuss

\*  $\text{Sym}^k(\mathbb{C}^2) \otimes \mathbb{C}^2$ :  $t_z^{(k)} \otimes I + I \otimes t_z^{(1)} =: r_z^{\text{LHS}}$   
 $|k, s\rangle \otimes |b\rangle \rightsquigarrow$  eigenvalue  $s \pm (-1)^b$

\*  $\text{Sym}^{k+1}(\mathbb{C}^2) \oplus \text{Sym}^{k-1}(\mathbb{C}^2)$ :  $t_z^{(k+1)} \oplus t_z^{(k-1)} =: r_z^{\text{RHS}}$   
 $|k+1, s'\rangle, |k-1, s'\rangle \rightsquigarrow$  eigenvalue  $s'$

$$J(|k, s' - (-1)^b\rangle \otimes |b\rangle) = \sum_{P=\pm 1} U(k, s)_{P, b} |k+p, s'\rangle$$

\* rewrite:

2x2 unitary matrix  $\infty$

(also possible for  $s' = \pm(k+1)$ )

$$J(|k, s\rangle \otimes |b\rangle) = \sum_{P=\pm 1} U(k, s + (-1)^b)_{P, b} |k+p, s + (-1)^b\rangle$$

\* How to compute!

$$J(|k, k\rangle \otimes |0\rangle) = |k+1, k+1\rangle \xrightarrow{r_{k-}} |k+1, k-1\rangle$$

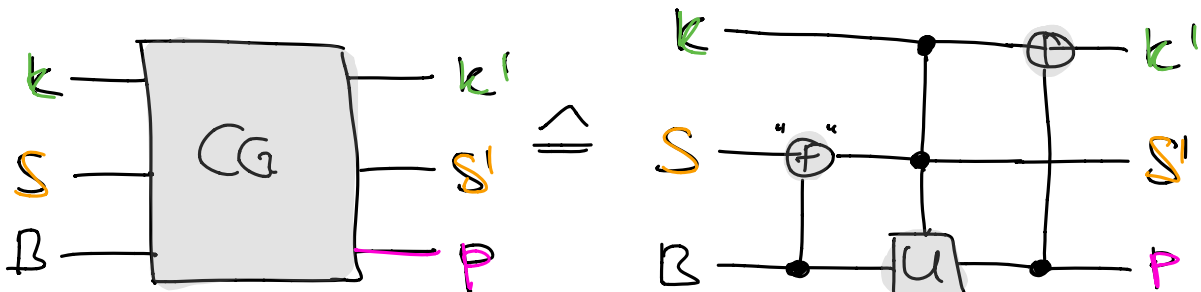
$$\xrightarrow{\pm} |k-1, k-1\rangle$$

& continue via  $r_{k-}$

Result:

$$\langle k', s' | J_k(|k, s\rangle \otimes |b\rangle)$$

$$|k\rangle \otimes |s\rangle \otimes |b\rangle \rightarrow \sum_P |k+p\rangle \otimes |s + (-1)^b\rangle \otimes U(k, s + (-1)^b)_{P, b} |P\rangle$$



Rest of today: Q-entropy & a new task...

# Entropy & Entropy inequalities

Shannon entropy:  $H(\{p_i\}_{i=1}^d) := - \sum_{i=1}^d p_i \log p_i$   
 $0 \cdot \log 0 := -1$

\*  $H(\{p_i\}) = h(p)$  ✓

Von Neumann entropy:  $\rho$  on  $\mathbb{C}^d$

$$S(\rho) := -\text{tr}[\rho \cdot \log \rho]$$

\*  $S(\rho) = H(\{p_i\})$  if  $\{p_i\}$  eigenvalues (with multiplicity!)

\*  $0 \leq S(\rho) \leq \log d$ ,  $= \log d$  iff  $\rho = \frac{I}{d}$   
 $\uparrow$  Jensen:  $\sum_i p_i \log(\frac{1}{p_i}) \leq \log(\sum_i p_i \frac{1}{p_i})$

\* Optimal asympt. rate for Compression, state transfers also if  $d \geq 2$  (we only did qubits)

\* Asymptotic equipartition property (AEP):

$\forall \epsilon > 0 \exists$  "typical" projectors  $\{P_n\}$  on  $(\mathbb{C}^d)^{\otimes n}$

①  $\text{tr}[P_n \rho^{\otimes n}] \rightarrow 1$

②  $\text{rk}[P_n] \leq 2^{n(S(\rho) + \epsilon)}$

③ eigenvalues of  $P_n \rho^{\otimes n} P_n$  are  $2^{-n(S(\rho) \pm \epsilon)}$

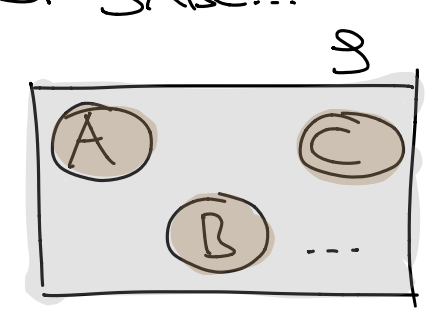
We/you proved this for  $d=2$ ! (③? PSET 6 :-)

SUMMARY:  $\rho^{\otimes n} \stackrel{\text{gentle meas.}}{\approx} P_n \rho^{\otimes n} P_n \leftarrow$  "uniform" on typ. subspace

# Entropies of subsystems:

\* **NOTATION:**  $S(A)_\rho := S(\rho_A)$  for  $\rho_{ABC} \dots$

↑ often omit



\*  $\rho_{AB}$  pure:  $S(AB) = 0$ ,  $S(A) = S(B) = S_E$

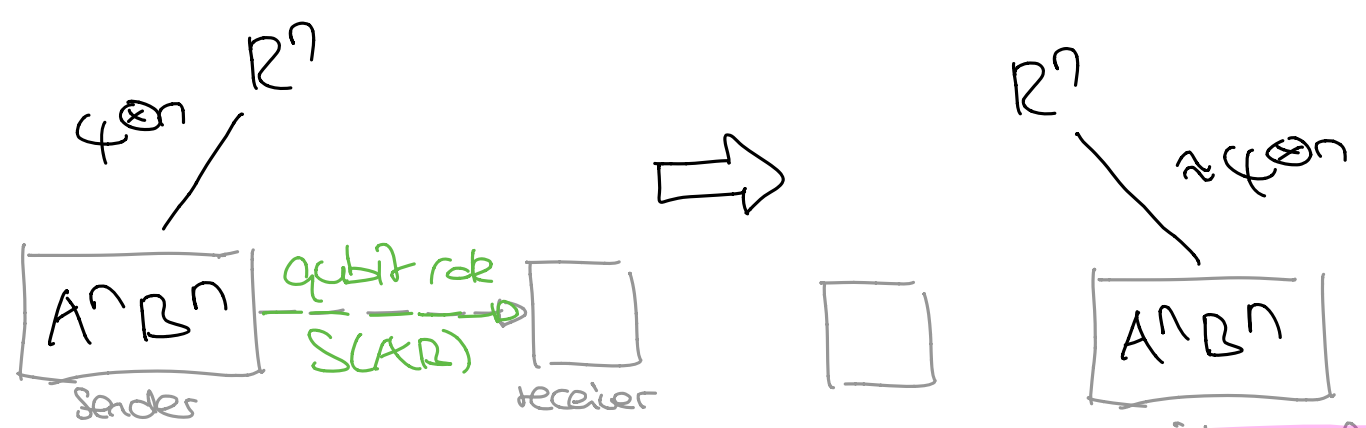
\*  $\rho_{AB} = \rho_A \otimes \rho_B$ :  $S(AB) = S(A) + S(B)$  ↕ compare

Pf:  $\rho_A \{p_i\}, \rho_B \{q_j\} \rightarrow \rho_{AB} \{p_i q_j\}$  &  
 $-\sum_{i,j} p_i q_j \log(p_i q_j) = -\sum_{i,j} p_i q_j \log(p_i) - \sum_{i,j} p_i q_j \log(q_j) \quad \square$

\* **Subadditivity:**  $S(AB)_\rho \leq S(A)_\rho + S(B)_\rho$

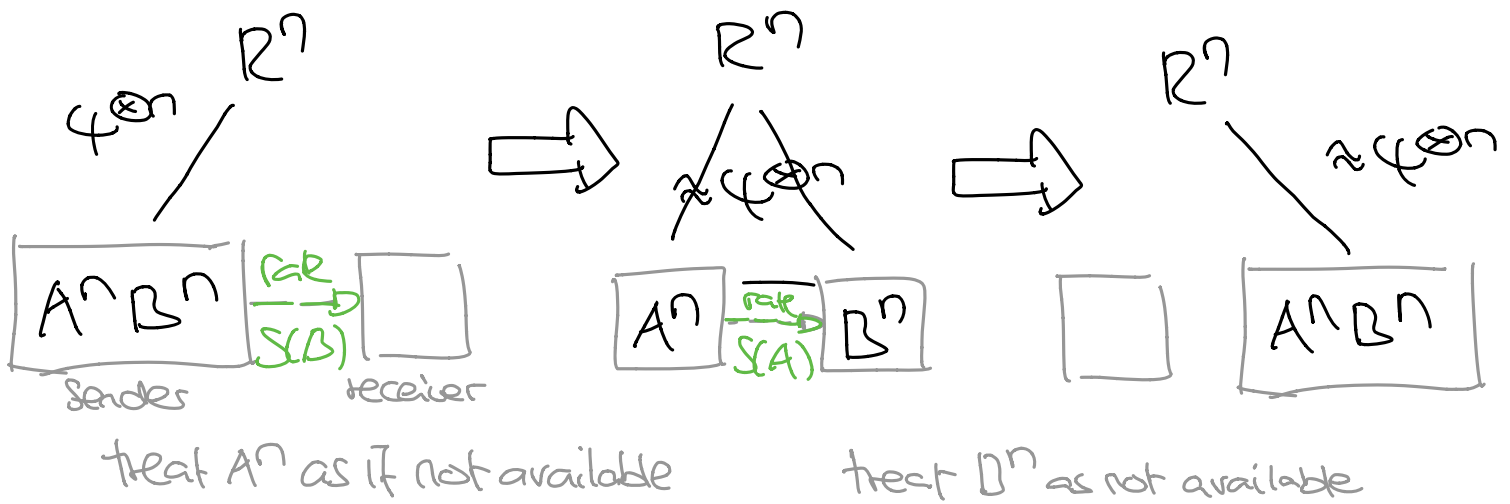
"Pf:" Let  $|\psi_{ABR}\rangle$  be a purification.

Recall **q. state transp.**



•  $S(AB)$  is optimal qubit rate ← we did not fully prove this (otherwise completely rigorous)

• ...but  $S(A)+S(B)$  is achievable!



Alternative analysis proof:  
Klein's inequality (□)

\* Araki-Lieb inequality:

$$S(AB) \geq |S(A) - S(B)|$$

Pf:  $S_{AB} \sim (\Psi_{ABR})$

$$S(AB) = S(C) \stackrel{SA}{\geq} S(BC) - S(B) = S(A) - S(B) \text{ etc. } \square$$

Mutual information:  $S_{AB}$  on  $\mathbb{C}^{d_A} \otimes \mathbb{C}^{d_B}$ :

$$I(A:B)_g := S(A)_g + S(B)_g - S(AB)_g$$

$\hat{=}$  info that we lose when we treat  $A+B$  independently

$$* I(A:B) \stackrel{SSA}{\geq} 0, = 0 \text{ iff } S_{AB} = S_A \otimes S_B$$

$$* S_{AB} \text{ pure: } \frac{1}{2} I(A:B) = S(A) = S(B) = S_E$$

\*  $I(A:B) \stackrel{\text{Araki-Lieb}}{\leq} 2 \cdot S(A) \leq 2 \cdot \log d_A$

likewise for B

\*  $I(A:B) > S(A) \implies S_{AB}$  is entangled !

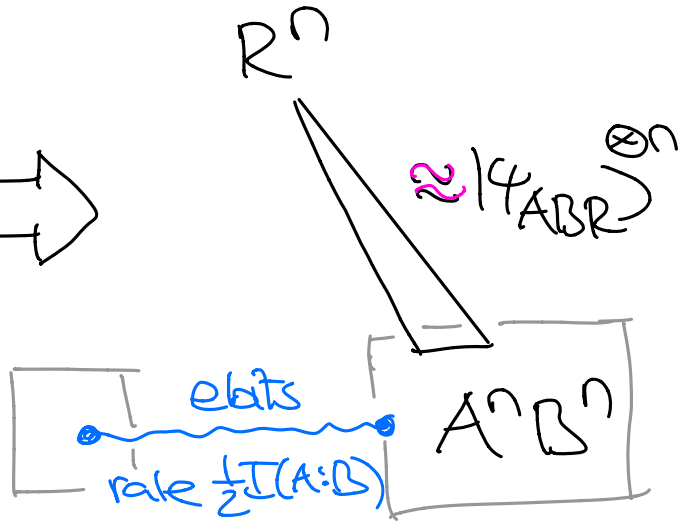
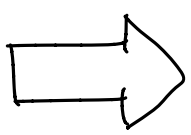
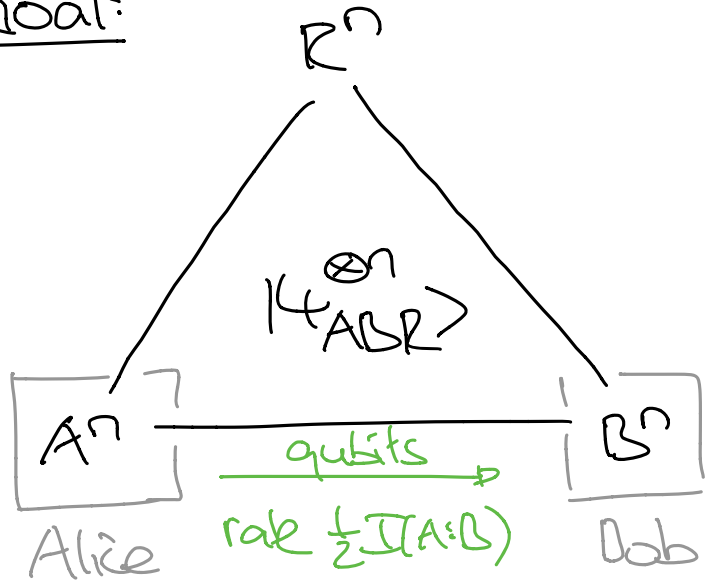
e.g.  $|\Phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$  :  $I(A:B) = 1 + 1 - 0 = 2$

$S_{AB} = \frac{1}{2}(|00\rangle\langle 00| + |11\rangle\langle 11|)$  :  $I(A:B) = 1 + 1 - 1 = 1$

### Quantum State merging

a.k.a. coherent q state merging, fully q. Stepan-Wolf, ...

Goal:



- \* For Comparison: q. state transfer needs qubit rate  $S(A)$ : yields no ebits
- \* If **no B**: Q. state transfer  $\rightarrow S(A) = \frac{1}{2} I(A:R)$ , optimal!
- \* If **no R**: Need not send anything! Bob can just prepare state locally  
 ↳ Can hope to extract entanglement...

RESULT (tomorrow):  $\exists$  protocol s.th. given  
 $I(A:B)$  on

① qubit rate  $\frac{1}{2} I(A:R)$  suffices

② get ebits at rate  $\frac{1}{2} I(A:B)$  for free

generalizes many QIT protocols!  $\nabla$