

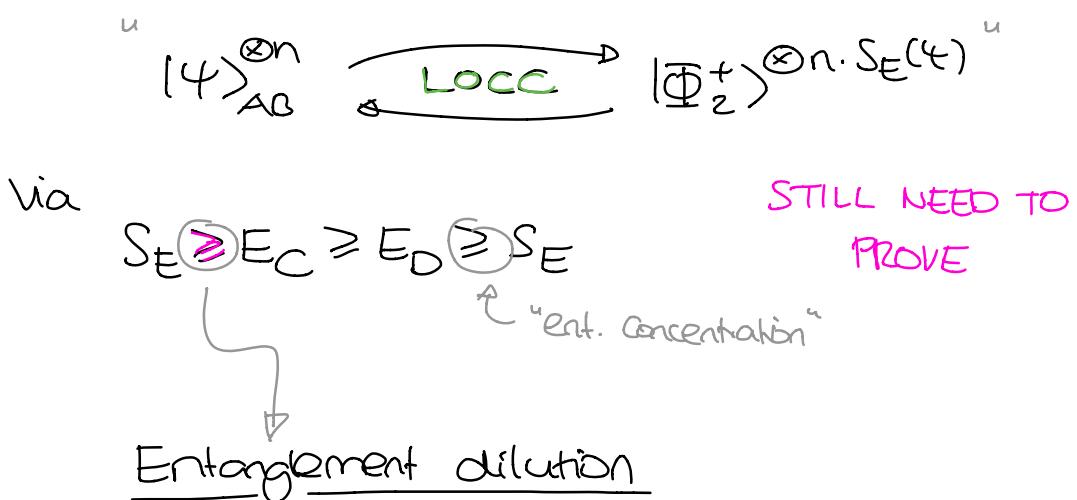
Last time: Entanglement entropy of $|4\rangle_{AB}^{\otimes n}$:

$$S_E(4) := S(g_A) = S(g_B)$$

① (Optimal) quantum compression rate

OPERATIONAL
INTERPRETATION

② $S_E = E_C = E_D$:



$$|\Phi^+\rangle_2^{\otimes R \cdot n} \xrightarrow{\text{LOCC}} \approx |\Psi\rangle_{AB}^{\otimes n}$$

Idea:

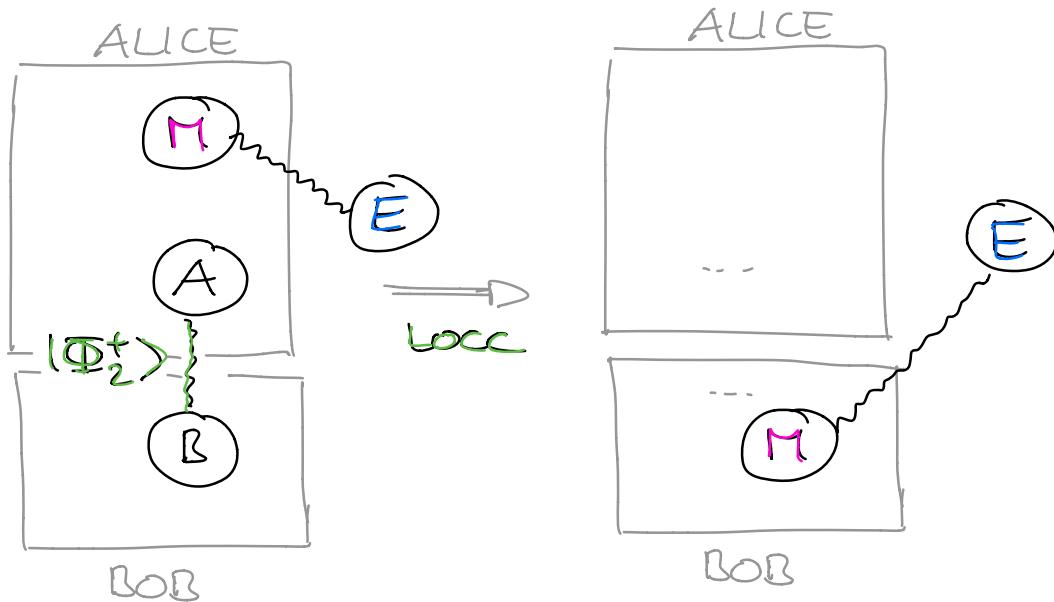
- Alice can prepare $|\Psi\rangle_{AB}^{\otimes n}$ in her laboratory.
- Q. compression allows to transfer B-systems to Bob by sending $\approx n \cdot (S_E + \delta)$ qubits

Q: Can we instead use ebits and LOCC?

Quantum Teleportation

Alice and Bob share ebit $| \Phi_2^+ \rangle_{AB}$

Goal: Send qubit M of Alice over to Bob.



$$|\psi\rangle_{EM} \otimes |\Phi_2^+\rangle_{AB}$$

Alice Bob

$$\dots \otimes |\psi\rangle_{EM}$$

Bob

(Entanglement of M should be preserved.)

- No cloning: Can only succeed if Alice learns nothing about state of M (no info remains in her lab)
- OTOH: Alice should couple A & M.

↳ measure AM in basis of max entangled states:

$$|\phi_0\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle) = (I \otimes I) |\Phi_2^+\rangle$$

$$|\phi_1\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle) = (I \otimes Z) |\Phi_2^+\rangle$$

$$|\phi_2\rangle = \frac{1}{\sqrt{2}}(|10\rangle + |01\rangle) = (I \otimes X) |\Phi_2^+\rangle$$

$$|\phi_3\rangle = \frac{1}{\sqrt{2}}(|10\rangle - |01\rangle) = (I \otimes XZ) |\Phi_2^+\rangle$$

i.e. $|\phi_k\rangle_{AM} = (I \otimes U_k) |\Phi_2^+\rangle_{AM}$

↳ proj. measurement $P_{AM,k} = |\phi_k\rangle \langle \phi_k|_{AM}$

$\Pr(\text{outcome } k)$

$$= \text{tr}[P_{AM,k} \cdot \text{tr}_E [|\psi\rangle \langle \psi|_{EM} \otimes |\Phi_2^+\rangle \langle \Phi_2^+|_{AB}]]$$

$$= \text{tr}[|\phi_k\rangle \langle \phi_k|_{AM} \cdot (\text{tr}_E [|\psi\rangle \langle \psi|_{EM}] \otimes \frac{I_A}{2})]$$

$$= \frac{1}{2} \text{tr}[\text{tr}_A [|\phi_k\rangle \langle \phi_k|_{AM}] \cdot \text{tr}_E [|\psi\rangle \langle \psi|_{EM}]]$$

$$= \frac{1}{2} \text{tr}[\frac{I_M}{2} \text{tr}_E [|\psi\rangle \langle \psi|_{EM}]] = \frac{1}{4}.$$

COMPLETELY UNINFORMATIVE ☺

Post-measurement state?

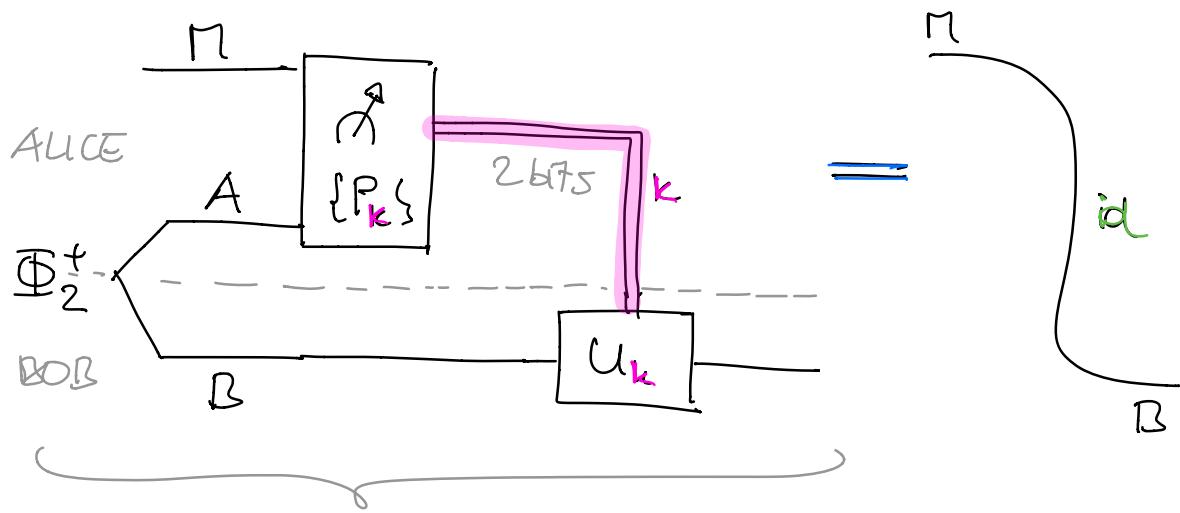
$$\begin{aligned} & 2 (|\phi_k\rangle \langle \phi_k|_{AM} \otimes I_E) (|\psi\rangle_{EM} \otimes |\Phi_2^+\rangle_{AB}) \\ &= 2 \underbrace{[I_E \otimes (|\Phi_2^+\rangle \langle \Phi_2^+|_{AM} \otimes I_B) (I_M \otimes |\Phi_2^+\rangle \langle \Phi_2^+|_{AB})]}_{=?} \underbrace{[I_E \otimes U_{M,k}^\dagger] |\psi\rangle_{EM}}_{=} \end{aligned}$$

$$\begin{aligned}
 & \downarrow \\
 & (\langle \Phi_2^+ |_{AB} \otimes I_B) (I_n \otimes |\Phi_2^+\rangle_{AB}) \\
 = & \frac{1}{2} \sum_{x,y} (\langle x|_A \otimes \langle x|_n \otimes I_B) (|y\rangle_A \otimes I_n \otimes |y\rangle_B) \\
 = & \frac{1}{2} \left[\sum_x |x\rangle_B \langle x|_n \right]
 \end{aligned}$$

IDENTITY FROM n TO B

$$\Rightarrow ? = (I_E \otimes U_{B_1}^+) |4\rangle_{EB}$$

Alice sends over k & Bob applies U_{B_1k} : SUCCESS



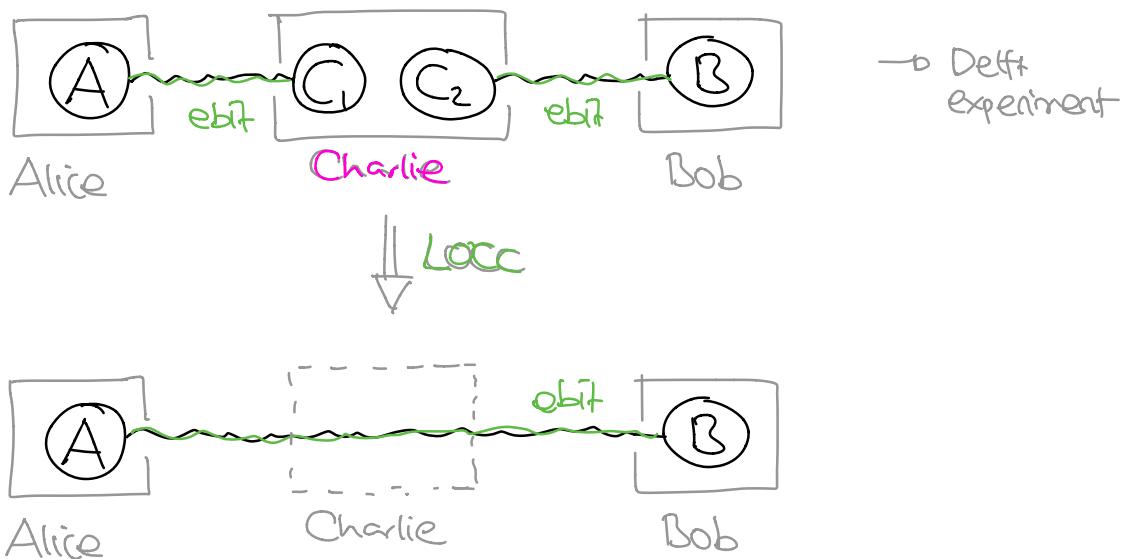
- LOCC! zero error, non-asymptotic
- Composable: n ebits $\xrightarrow{\text{LOCC}}$ n qubits

↳ Can convert any compression protocol into entanglement dilution protocol at same rate

$$\Rightarrow E_S = R_{\text{opt Compr.}}^{\text{opt}} - E_C = E_D$$

everything we wanted to prove

Entanglement swapping:



- Similarly with many intermediate "relays"!

Resource inequalities

Teleportation:

$$\text{ebit} + 2 \underbrace{[c \rightarrow c]}_{\substack{1 \text{ bit of class-} \\ \text{Commun.}}} \geq \underbrace{[q \rightarrow q]}_{\substack{1 \text{ qubit of quantum} \\ \text{communication}}}$$

Local Operations

What else can we do?

- $[q \rightarrow q] \geq \text{ebit}$ ← Alice prepares ebit, sends over half of it

BUT: $\text{ebit} \cancel{\geq} [q \rightarrow q]$ CANNOT COMMUNICATE
USING ENTANGLEMENT ALONE

- $[q \rightarrow q] \geq [c \rightarrow c]$ ← Alice encodes x by $\{x\}$, Bob measures $\{xx\}$

BUT: $[q \rightarrow q] \cancel{\geq} 2[c \rightarrow c]$ "HOLEVO BOUND"

Superdense coding:

(not quite) converse

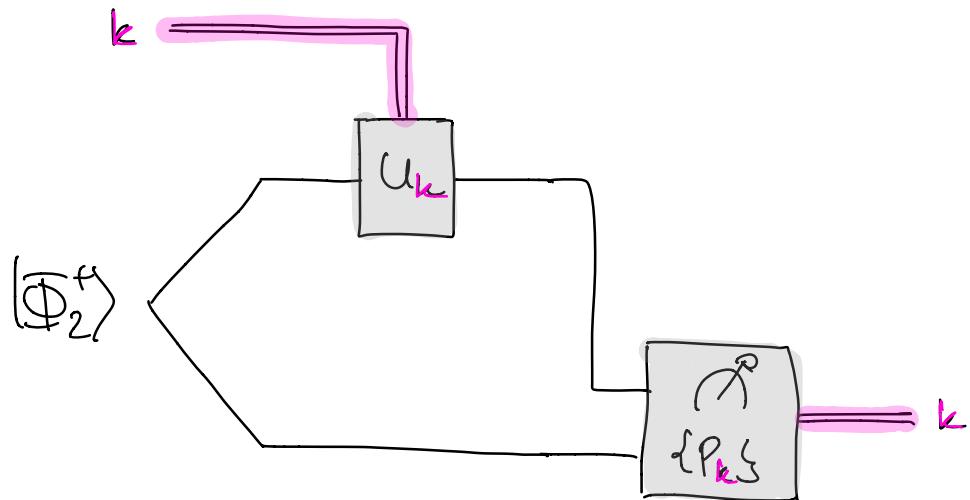
$$[q \rightarrow q] + \text{ebit} \geq 2[c \rightarrow c]$$

of teleportation

$$\Rightarrow [q \rightarrow q] = 2[c \rightarrow c] \bmod \text{ebit}$$

Protocol: • Alice and Bob share $(\tilde{\Phi}_2^+)_A B$

- Alice applies some U_k to A & sends it over to Bob
- Bob measures $\{P_k\}$



Resource inequalities are very useful to relate QIP protocols in their strength.